

Integrated Circuits Group

LH28F320BJHE-PBTL90

Flash Memory 32M (2M x 16/4M × 8)

(Model No.: LHF32J08)

Spec No.: EL12X044 Issue Date: October 16, 2000

SHARP

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PAGE

CONTENTS

5

6

PAGE

1 INTRODUCTION.31.1 Features31.2 Product Overview.31.3 Product Description41.3.1 Package Pinout41.3.2 Block Organization.4

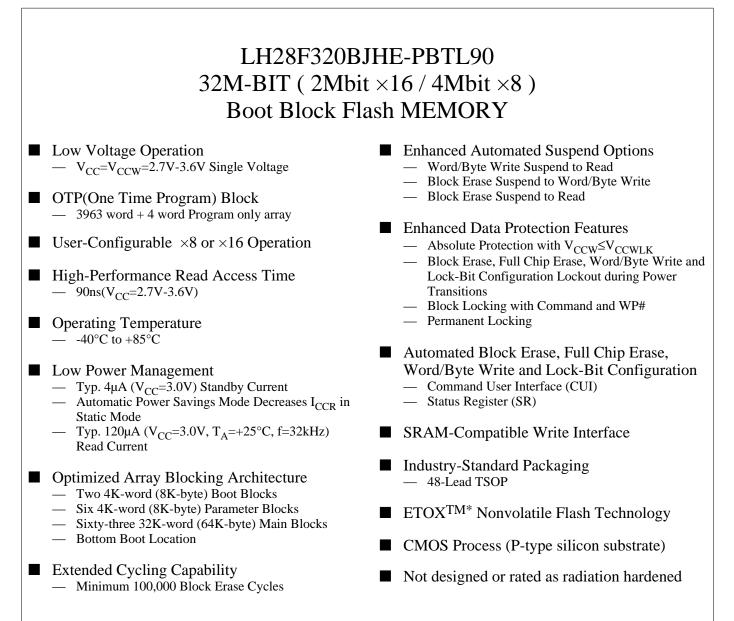
2 PRINCIPLES OF OPERATION	8
2.1 Data Protection	8

3 BUS OPERATION	9
3.1 Read	9
3.2 Output Disable	9
3.3 Standby	9
3.4 Reset	9
3.5 Read Identifier Codes	10
3.6 OTP(One Time Program) Block	10
3.7 Write	11

4 COMMAND DEFINITIONS114.1 Read Array Command134.2 Read Identifier Codes Command134.3 Read Status Register Command134.4 Clear Status Register Command134.5 Block Erase Command144.6 Full Chip Erase Command144.7 Word/Byte Write Command144.8 Block Erase Suspend Command154.9 Word/Byte Write Suspend Command154.10 Set Block and Permanent Lock-Bit Command164.11 Clear Block Lock-Bits Command174.13 Block Locking by the WP#17

DESIGN CONSIDERATIONS
5.1 Three-Line Output Control
5.2 RY/BY# and WSM Polling27
5.3 Power Supply Decoupling
5.4 V _{CCW} Trace on Printed Circuit Boards 27
5.5 V _{CC} , V _{CCW} , RP# Transitions
5.6 Power-Up/Down Protection
5.7 Power Dissipation
5.8 Data Protection Method
ELECTRICAL SPECIFICATIONS
6.1 Absolute Maximum Ratings
6.2 Operating Conditions
6.2.1 Capacitance
6.2.2 AC Input/Output Test Conditions 30
6.2.3 DC Characteristics 31
6.2.4 AC Characteristics - Read-Only Operations 33
6.2.5 AC Characteristics - Write Operations
6.2.6 Alternative CE#-Controlled Writes
6.2.7 Reset Operations 40
6.2.8 Block Erase, Full Chip Erase, Word/Byte Write and
Lock-Bit Configuration Performance 41





SHARP's LH28F320BJHE-PBTL90 Flash memory is a high-density, low-cost, nonvolatile, read/write storage solution for a wide range of applications.

LH28F320BJHE-PBTL90 can operate at V_{CC} =2.7V-3.6V and V_{CCW} =2.7V-3.6V or 11.7V-12.3V. Its low voltage operation capability realize battery life and suits for cellular phone application.

Its Boot, Parameter and Main-blocked architecture, low voltage and extended cycling provide for highly flexible component suitable for portable terminals and personal computers. Its enhanced suspend capabilities provide for an ideal solution for code + data storage applications.

For secure code storage applications, such as networking, where code is either directly executed out of flash or downloaded to DRAM, the LH28F320BJHE-PBTL90 offers four levels of protection: absolute protection with $V_{CCW} \leq V_{CCWLK}$, selective hardware block locking or flexible software block locking. These alternatives give designers ultimate control of their code security needs.

The LH28F320BJHE-PBTL90 is manufactured on SHARP's 0.25 μ m ETOX^{TM*} process technology. It come in industry-standard package: the 48-lead TSOP, ideal for board constrained applications.

*ETOX is a trademark of Intel Corporation.

1 INTRODUCTION

This datasheet contains LH28F320BJHE-PBTL90 specifications. Section 1 provides a flash memory overview. Sections 2, 3, 4 and 5 describe the memory organization and functionality. Section 6 covers electrical specifications.

1.1 Features

Key enhancements of LH28F320BJHE-PBTL90 boot block Flash memory are:

- •Single low voltage operation
- •Low power consumption
- •Enhanced Suspend Capabilities
- •Boot Block Architecture

Please note following:

1.2 Product Overview

The LH28F320BJHE-PBTL90 is a high-performance 32M-bit Boot Block Flash memory organized as 2M-word of 16 bits or 4M-byte of 8 bits. The 2M-word/4M-byte of data is arranged in two 4K-word/8K-byte boot blocks, six 4K-word/8K-byte parameter blocks and sixty-three 32K-word/64K-byte main blocks which are individually erasable, lockable and unlockable in-system. The memory map is shown in Figure 3.

The dedicated V_{CCW} pin gives complete data protection when $V_{CCW} \leq V_{CCWLK}$.

A Command User Interface (CUI) serves as the interface between the system processor and internal operation of the device. A valid command sequence written to the CUI initiates device automation. An internal Write State Machine (WSM) automatically executes the algorithms and timings necessary for block erase, full chip erase, word/byte write and lock-bit configuration operations. A block erase operation erases one of the device's 32K-word/64K-byte blocks typically within 1.2s (3V V_{CC} , 3V V_{CCW}), 4K-word/8K-byte blocks typically within 0.6s (3V V_{CC} , 3V V_{CCW}) independent of other blocks. Each block can be independently erased minimum 100,000 times. Block erase suspend mode allows system software to suspend block erase to read or write data from any other block.

Writing memory data is performed in word/byte increments of the device's 32K-word blocks typically within 33 μ s (3V V_{CC}, 3V V_{CCW}), 64K-byte blocks typically within 31 μ s (3V V_{CC}, 3V V_{CCW}), 4K-word blocks typically within 36 μ s (3V V_{CC}, 3V V_{CCW}), 8K-byte blocks typically within 32 μ s (3V V_{CC}, 3V V_{CCW}), 8K-byte blocks typically within 32 μ s (3V V_{CC}, 3V V_{CCW}). Word/byte write suspend mode enables the system to read data or execute code from any other flash memory array location.

Individual block locking uses a combination of bits, seventy-one block lock-bits, a permanent lock-bit and WP# pin, to lock and unlock blocks. Block lock-bits gate block erase, full chip erase and word/byte write operations, while the permanent lock-bit gates block lockbit modification and locked block alternation. Lock-bit configuration operations (Set Block Lock-Bit, Set Permanent Lock-Bit and Clear Block Lock-Bits commands) set and cleared lock-bits.

The status register indicates when the WSM's block erase, full chip erase, word/byte write or lock-bit configuration operation is finished.

The RY/BY# output gives an additional indicator of WSM activity by providing both a hardware signal of status (versus software polling) and status masking (interrupt masking for background block erase, for example). Status polling using RY/BY# minimizes both CPU overhead and system power consumption. When low, RY/BY# indicates that the WSM is performing a block erase, full chip erase, word/byte write or lock-bit configuration. RY/BY#-high Z indicates that the WSM is ready for a new command, block erase is suspended (and word/byte write is inactive), word/byte write is suspended, or the device is in reset mode.

The access time is 90ns (t_{AVQV}) over the operating temperature range (-40°C to +85°C) and V_{CC} supply voltage range of 2.7V-3.6V.

The Automatic Power Savings (APS) feature substantially reduces active current when the device is in static mode (addresses not switching). In APS mode, the typical I_{CCR} current is 4µA (CMOS) at 3.0V V_{CC} .

When CE# and RP# pins are at V_{CC} , the I_{CC} CMOS standby mode is enabled. When the RP# pin is at GND, reset mode is enabled which minimizes power consumption and provides write protection. A reset time (t_{PHQV}) is required from RP# switching high until outputs are valid. Likewise, the device has a wake time (t_{PHEL}) from RP#-high until writes to the CUI are recognized. With RP# at GND, the WSM is reset and the status register is cleared.

Please do not execute reprogramming "0" for the bit which has already been programed "0". Overwrite operation may generate unerasable bit. In case of reprogramming "0" to the data which has been programed "1".

•Program "0" for the bit in which you want to change data from "1" to "0".

•Program "1" for the bit which has already been programmed "0".

For example, changing data from "10111101" to "10111100" requires "1111110" programming.

1.3 Product Description

1.3.1 Package Pinout

LH28F320BJHE-PBTL90 Boot Block Flash memory is available in 48-lead TSOP package (see Figure 2).

1.3.2 Block Organization

This product features an asymmetrically-blocked architecture providing system memory integration. Each erase block can be erased independently of the others up to 100,000 times. For the address locations of the blocks, see the memory map in Figure 3.

Boot Blocks: The boot block is intended to replace a dedicated boot PROM in a microprocessor or microcontroller-based system. This boot block 4K words (4,096words) features hardware controllable write-protection to protect the crucial microprocessor boot code from accidental modification. The protection of the boot block is controlled using a combination of the V_{CCW}, RP#, WP# pins and block lock-bit.

Parameter Blocks: The boot block architecture includes parameter blocks to facilitate storage of frequently update small parameters that would normally require an EEPROM. By using software techniques, the word-rewrite functionality of EEPROMs can be emulated. Each boot block component contains six parameter blocks of 4K words (4,096 words) each. The protection of the parameter block is controlled using a combination of the V_{CCW}, RP# and block lock-bit.

Main Blocks: The reminder is divided into main blocks for data or code storage. Each 32M-bit device contains sixty-three 32K words (32,768 words) blocks. The protection of the main block is controlled using a combination of the V_{CCW} , RP# and block lock-bit.

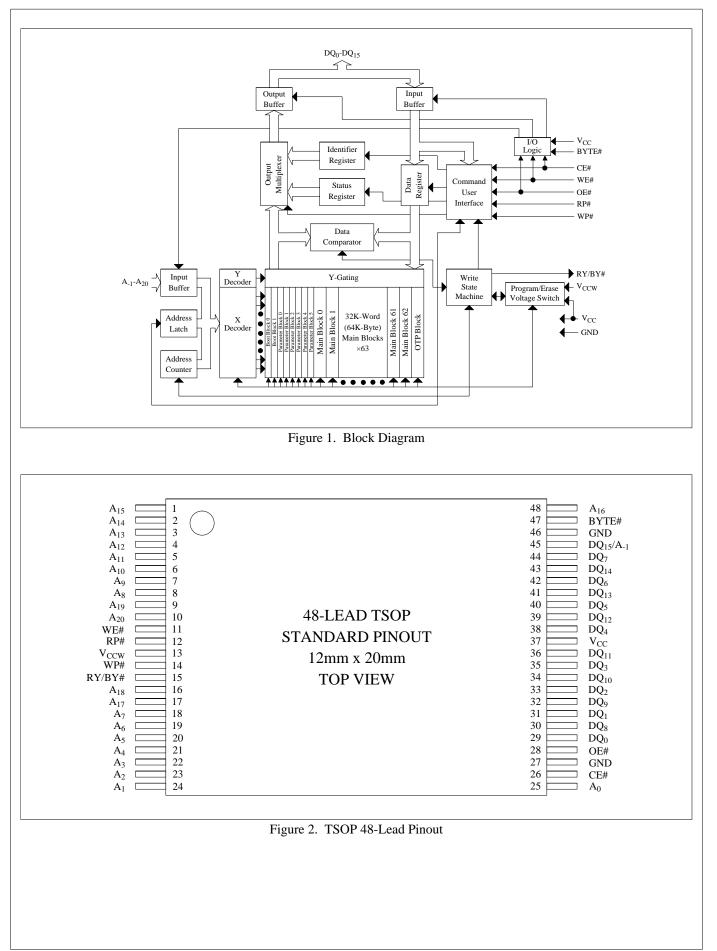


		Table 1. Pin Descriptions
Symbol	Туре	Name and Function
A ₋₁ A ₀ -A ₂₀	INPUT	ADDRESS INPUTS: Inputs for addresses during read and write operations. Addresses are internally latched during a write cycle. A ₋₁ : Lower address input while BYTE# is V_{IL} . A ₋₁ pin changes DQ ₁₅ pin while BYTE# is V_{IH} . A ₁₅ -A ₂₀ : Main Block Address. A ₁₂ -A ₂₀ : Boot and Parameter Block Address.
DQ ₀ -DQ ₁₅	INPUT/ OUTPUT	DATA INPUT/OUTPUTS: Inputs data and commands during CUI write cycles; outputs data during memory array, status register and identifier code read cycles. Data pins float to high-impedance when the chip is deselected or outputs are disabled. Data is internally latched during a write cycle. DQ_8 - DQ_{15} pins are not used while byte mode (BYTE#= V_{IL}). Then, DQ_{15} pin changes A_{-1} address input.
CE#	INPUT	CHIP ENABLE: Activates the device's control logic, input buffers, decoders and sense amplifiers. CE#-high deselects the device and reduces power consumption to standby levels.
RP#	INPUT	RESET: Resets the device internal automation. RP#-high enables normal operation. When driven low, RP# inhibits write operations which provides data protection during power transitions. Exit from reset mode sets the device to read array mode. RP# must be V_{IL} during power-up.
OE#	INPUT	OUTPUT ENABLE: Gates the device's outputs during a read cycle.
WE#	INPUT	WRITE ENABLE: Controls writes to the CUI and array blocks. Addresses and data are latched on the rising edge of the WE# pulse.
WP#	INPUT	WRITE PROTECT: When WP# is V_{IL} , boot blocks cannot be written or erased. When WP# is V_{IH} , locked boot blocks can not be written or erased. WP# is not affected parameter and main blocks.
BYTE#	INPUT	BYTE ENABLE: BYTE# V_{IL} places device in byte mode (×8). All data is then input or output on DQ_{0-7} , and DQ_{8-15} float. BYTE# V_{IH} places the device in word mode (×16), and turns off the A_{-1} input buffer.
RY/BY#	OPEN DRAIN OUTPUT	READY/BUSY#: Indicates the status of the internal WSM. When low, the WSM is performing an internal operation (block erase, full chip erase, word/byte write or lock-bit configuration). RY/BY#-high Z indicates that the WSM is ready for new commands, block erase is suspended, and word/byte write is inactive, word/byte write is suspended, or the device is in reset mode.
V _{CCW}	SUPPLY	BLOCK ERASE, FULL CHIP ERASE, WORD/BYTE WRITE OR LOCK-BIT CONFIGURATION POWER SUPPLY: For erasing array blocks, writing words/bytes or configuring lock-bits. With $V_{CCW} \leq V_{CCWLK}$, memory contents cannot be altered. Block erase, full chip erase, word/byte write and lock-bit configuration with an invalid V_{CCW} (see 6.2.3 DC Characteristics) produce spurious results and should not be attempted. Applying 12V±0.3V to V_{CCW} during erase/write can only be done for a maximum of 1000 cycles on each block. V_{CCW} may be connected to 12V±0.3V for a total of 80 hours maximum.
V _{CC}	SUPPLY	DEVICE POWER SUPPLY: Do not float any power pins. With $V_{CC} \leq V_{LKO}$, all write attempts to the flash memory are inhibited. Device operations at invalid V_{CC} voltage (see 6.2.3 DC Characteristics) produce spurious results and should not be attempted.
GND	SUPPLY	GROUND: Do not float any ground pins.

[A ₂₀ -A ₀]			[A ₂₀ -A ₋₁]
0FFFFF 0F8000	32KW/64KB Main Block	30	1FFFFF
0F7FFF	32KW/64KB Main Block	29	1F0000 1EFFFF
0F0000 0EFFFF	32KW/64KB Main Block	28	1E0000 1DFFFF
0E8000 0E7FFF	32KW/64KB Main Block	27	1D0000 1CFFFF
0E0000 0DFFFF	32KW/64KB Main Block	26	1C0000 1BFFFF
0D8000 0D7FFF	32KW/64KB Main Block	25	1B0000 1AFFFF
0D0000 0CFFFF	32KW/64KB Main Block	23	1A0000 19FFFF
0C8000 0C7FFF	32KW/64KB Main Block	24	190000 18FFFF
0C0000 0BFFFF			180000 17FFFF
0B8000 0B7FFF	32KW/64KB Main Block	22	170000 16FFFF
0B0000 0AFFFF	32KW/64KB Main Block	21	160000 15FFFF
0A8000 0A7FFF	32KW/64KB Main Block	20	150000 14FFFF
0A0000	32KW/64KB Main Block	19	140000
09FFFF 098000	32KW/64KB Main Block	18	13FFFF 130000
097FFF 090000	32KW/64KB Main Block	17	12FFFF 120000
08FFFF 088000	32KW/64KB Main Block	16	11FFFF 110000
087FFF 080000	32KW/64KB Main Block	15	10FFFF 100000
07FFFF 078000	32KW/64KB Main Block	14	0FFFFF 0F0000
077FFF 070000	32KW/64KB Main Block	13	0EFFFF 0E0000
06FFFF 068000	32KW/64KB Main Block	12	0DFFFF 0D0000
068000 067FFF 060000	32KW/64KB Main Block	11	0CFFFF
05FFFF	32KW/64KB Main Block	10	0C0000 0BFFFF
058000 057FFF	32KW/64KB Main Block	9	0B0000 0AFFFF
050000 04FFFF	32KW/64KB Main Block	8	0A0000 09FFFF
048000 047FFF	32KW/64KB Main Block	7	090000 08FFFF
040000 03FFFF	32KW/64KB Main Block	6	080000 07FFFF
038000 037FFF	32KW/64KB Main Block	5	070000 06FFFF
030000 02FFFF	32KW/64KB Main Block	4	060000 05FFFF
028000 027FFF			050000 04FFFF
020000 01FFFF	32KW/64KB Main Block	3	040000 03FFFF
018000 017FFF	32KW/64KB Main Block	2	030000 02FFFF
010000 00FFFF	32KW/64KB Main Block	1	020000 01FFFF
008000	32KW/64KB Main Block	0	010000
007FFF 007000	4KW/8KB Parameter Block	5	00FFFF 00E000
006FFF 006000	4KW/8KB Parameter Block	4	00DFFF 00C000
005FFF 005000	4KW/8KB Parameter Block	3	00BFFF 00A000
004FFF 004000	4KW/8KB Parameter Block	2	009FFF 008000
003FFF 003000	4KW/8KB Parameter Block	1	007FFF 006000
002FFF 002000	4KW/8KB Parameter Block	0	005FFF 004000
001FFF	4KW/8KB Boot Block	1	003FFF
001000 000FFF	4KW/8KB Boot Block	0	002000 001FFF
000000		2	000000

1]	$[A_{20}-A_0]$		
	1FFFFF	32KW/64KB Main Block	62
7	1F8000 1F7FFF	32KW/64KB Main Block	61
7	1F0000 1EFFFF	32KW/64KB Main Block	60
7	1E8000 1E7FFF	32KW/64KB Main Block	59
7	1E0000 1DFFFF	32KW/64KB Main Block	58
7	1D8000 1D7FFF	32KW/64KB Main Block	
	1D0000 1CFFFF		57
	1C8000 1C7FFF	32KW/64KB Main Block	56
	1C0000 1BFFFF	32KW/64KB Main Block	55
	1B8000	32KW/64KB Main Block	54
	1B7FFF 1B0000	32KW/64KB Main Block	53
	1AFFFF 1A8000	32KW/64KB Main Block	52
	1A7FFF 1A0000	32KW/64KB Main Block	51
	19FFFF 198000	32KW/64KB Main Block	50
	197FFF 190000	32KW/64KB Main Block	49
	18FFFF 188000	32KW/64KB Main Block	48
	187FFF	32KW/64KB Main Block	47
	180000 17FFFF	32KW/64KB Main Block	46
	178000 177FFF	32KW/64KB Main Block	45
	170000 16FFFF	32KW/64KB Main Block	44
	168000 167FFF	32KW/64KB Main Block	43
	160000 15FFFF		-
	158000 157FFF	32KW/64KB Main Block	42
	150000 14FFFF	32KW/64KB Main Block	41
	148000 147FFF	32KW/64KB Main Block	40
	140000 13FFFF	32KW/64KB Main Block	39
	138000	32KW/64KB Main Block	38
	137FFF 130000	32KW/64KB Main Block	37
	12FFFF 128000	32KW/64KB Main Block	36
	127FFF 120000	32KW/64KB Main Block	35
	11FFFF 118000	32KW/64KB Main Block	34
	117FFF 110000	32KW/64KB Main Block	33
	10FFFF	32KW/64KB Main Block	32
	108000 107FFF	32KW/64KB Main Block	31

Figure 3. Memory Map

2 PRINCIPLES OF OPERATION

The LH28F320BJHE-PBTL90 flash memory includes an on-chip WSM to manage block erase, full chip erase, word/byte write and lock-bit configuration functions. It allows for: fixed power supplies during block erase, full chip erase, word/byte write and lock-bit configuration, and minimal processor overhead with RAM-like interface timings.

After initial device power-up or return from reset mode (see section 3 Bus Operations), the device defaults to read array mode. Manipulation of external memory control pins allow array read, standby and output disable operations.

Status register and identifier codes can be accessed through the CUI independent of the V_{CCW} voltage. High voltage on V_{CCW} enables successful block erase, full chip erase, word/byte write and lock-bit configurations. All functions associated with altering memory contents–block erase, full chip erase, word/byte write, lock-bit configuration, status and identifier codes–are accessed via the CUI and verified through the status register.

Commands are written using standard microprocessor write timings. The CUI contents serve as input to the WSM, which controls the block erase, full chip erase, word/byte write and lock-bit configuration. The internal algorithms are regulated by the WSM, including pulse repetition, internal verification and margining of data. Addresses and data are internally latched during write cycles. Writing the appropriate command outputs array data, accesses the identifier codes or outputs status register data. Interface software that initiates and polls progress of block erase, full chip erase, word/byte write and lock-bit configuration can be stored in any block. This code is copied to and executed from system RAM during flash memory updates. After successful completion, reads are again possible via the Read Array command. Block erase suspend allows system software to suspend a block erase to read/write data from/to blocks other than that which is suspend. Word/byte write suspend allows system software to suspend a word/byte write to read data from any other flash memory array location.

2.1 Data Protection

When $V_{CCW} \le V_{CCWLK}$, memory contents cannot be altered. The CUI, with two-step block erase, full chip erase, word/byte write or lock-bit configuration command sequences, provides protection from unwanted operations even when high voltage is applied to V_{CCW} . All write functions are disabled when V_{CC} is below the write lockout voltage V_{LKO} or when RP# is at V_{IL} . The device's block locking capability provides additional protection from inadvertent code or data alteration by gating block erase, full chip erase and word/byte write operations. Refer to Table 5 for write protection alternatives.

9

3 BUS OPERATION

The local CPU reads and writes flash memory in-system. All bus cycles to or from the flash memory conform to standard microprocessor bus cycles.

3.1 Read

Information can be read from any block, identifier codes or status register independent of the V_{CCW} voltage. RP# can be at V_{IH} .

The first task is to write the appropriate read mode command (Read Array, Read Identifier Codes or Read Status Register) to the CUI. Upon initial device power-up or after exit from reset mode, the device automatically resets to read array mode. Six control pins dictate the data flow in and out of the component: CE#, OE#, BYTE#, WE#, RP# and WP#. CE# and OE# must be driven active to obtain data at the outputs. CE# is the device selection control, and when active enables the selected memory device. OE# is the data output (DQ₀-DQ₁₅) control and when active drives the selected memory data onto the I/O bus. BYTE# is the device I/O interface mode control. WE# must be at V_{IL} or V_{IH}. Figure 16, 17 illustrates read cycle.

3.2 Output Disable

With OE# at a logic-high level (V_{IH}), the device outputs are disabled. Output pins (DQ_0 - DQ_{15}) are placed in a high-impedance state.

3.3 Standby

CE# at a logic-high level (V_{IH}) places the device in standby mode which substantially reduces device power consumption. DQ_0 - DQ_{15} outputs are placed in a high-impedance state independent of OE#. If deselected during block erase, full chip erase, word/byte write or lock-bit configuration, the device continues functioning, and consuming active power until the operation completes.

3.4 Reset

RP# at V_{II} initiates the reset mode.

In read modes, RP#-low deselects the memory, places output drivers in a high-impedance state and turns off all internal circuits. RP# must be held low for a minimum of 100ns. Time t_{PHQV} is required after return from reset mode until initial memory access outputs are valid. After this wake-up interval, normal operation is restored. The CUI is reset to read array mode and status register is set to 80H.

During block erase, full chip erase, word/byte write or lock-bit configuration modes, RP#-low will abort the operation. RY/BY# remains low until the reset operation is complete. Memory contents being altered are no longer valid; the data may be partially erased or written. Time t_{PHWL} is required after RP# goes to logic-high (V_{IH}) before another command can be written.

As with any automated device, it is important to assert RP# during system reset. When the system comes out of reset, it expects to read from the flash memory. Automated flash memories provide status information when accessed during block erase, full chip erase, word/byte write or lock-bit configuration modes. If a CPU reset occurs with no flash memory reset, proper CPU initialization may not occur because the flash memory may be providing status information instead of array data. SHARP's flash memories allow proper CPU initialization following a system reset through the use of the RP# input. In this application, RP# is controlled by the same RESET# signal that resets the system CPU.

3.5 Read Identifier Codes

The read identifier codes operation outputs the manufacturer code, device code, block lock configuration codes for each block and the permanent lock configuration code (see Figure 4). Using the manufacturer and device codes, the system CPU can automatically match the device with its proper algorithms. The block lock and permanent lock configuration codes identify locked and unlocked blocks and permanent lock-bit setting.

[A ₂₀ -A ₀]	Bottom Boot	[A ₂₀ -A ₋₁]
1FFFFF		3FFFFF
	Reserved for Future Implementation	
1F8003		3F0006
1F8002	Main Block 62 Lock Configuration Code	3F0005 3F0004
1F8001	Reserved for Future Implementation	3F0003
1F8000	Main Block 62	3F0000
1F7FFF	(Main Blocks 1 through 61)	3EFFFF
010000	(Main Diocks 1 unough 01)	020000
00FFFF		01FFFF
	Reserved for Future Implementation	
008003		001006
008002	Main Block 0 Lock Configuration Code	010005
008001	Reserved for Future Implementation	010003
008000	Main Block 0	010000
007FFF		00FFFF
	Reserved for Future Implementation	
007003		00E006
007002	Parameter Block 5 Lock Configuration Code	00E005 00E004
007001	Reserved for Future Implementation	00E003
007000	Parameter Block 5	00E000
006FFF	(Parameter Blocks 1 through 4)	00DFFI
003000	(Futuriotor Diocks Futiough 7)	006000
002FFF		005FFF
	Reserved for Future Implementation	
002003		004006
002002	Parameter Block 0 Lock Configuration Code	004005
002001	Reserved for Future Implementation	004003
002000	Parameter Block 0	004000
001FFF		003FFF
001003	Reserved for Future Implementation	002006
001003	Boot Block 1 Lock Configuration Code	002000
		002004
001001	Reserved for Future Implementation	
001000	Boot Block 1	002000
000FFF	OTP Block	001FFF
000080		000100
00007F 000004	Reserved for Future Implementation	0000FF 000008
000003	Permanent Lock Configuration Code	000007
000003	Boot Block 0 Lock Configuration Code	000006
000002	Device Code	000004
		000002
000000	Manufacturer Code Boot Block 0	000000

Figure 4. Device Identifier Code Memory Map

3.6 OTP(One Time Program) Block

The OTP block is a special block that can not be erased. The block is divided into two parts. One is a factory program area where a unique number can be written according to customer requirements in SHARP factory. This factory program area is "READ ONLY" (Already locked). The other is a customer program area that can be used by customers. This customer program area can be locked. After locking, this customer program area is protected permanently.

The OTP block is read in Configuration Read Mode by writing Read Identifier Codes command(90H). To return to Read Array Mode, write Read Array command(FFH).

The OTP block is programmed by writing OTP Program command(COH). First write OTP Program command and then write data with address to the device (See Figure 5). If OTP program is failed, SR.4(WORD/BYTE WRITE AND SET LOCK-BIT STATUS) bit is set to "1". And if this OTP block is locked, SR.1(DEVICE PROTECT STATUS) bit is set to "1" too.

The OTP block is also locked by writing OTP Program command(C0H). First write OTP Program command and then write data "FFFDH" with address "80H" to the device. Address "80H" of OTP block is OTP lock information. Bit 0 of address "80H" means factory program area lock status("1" is "NOT LOCKED", "0" is "LOCKED"). Bit 1 of address "80H" means customer program area lock status. The OTP lock information can not be cleared, after once it is set.

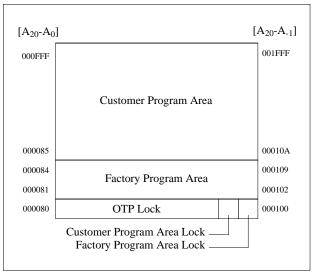


Figure 5. OTP Block Address Map

3.7 Write

Writing commands to the CUI enable reading of device data and identifier codes. They also control inspection and clearing of the status register. When V_{CC} =2.7V-3.6V and $V_{CCW} = V_{CCWH1/2}$, the CUI additionally controls block erase, full chip erase, word/byte write and lock-bit configuration.

The Block Erase command requires appropriate command data and an address within the block to be erased. The Full Chip Erase command requires appropriate command data and an address within the device. The Word/Byte Write command requires the command and address of the location to be written. Set Permanent and Block Lock-Bit commands require the command and address within the device (Permanent Lock) or block within the device (Block Lock) to be locked. The Clear Block Lock-Bits command requires the command and address within the device.

The CUI does not occupy an addressable memory location. It is written when WE# and CE# are active. The address and data needed to execute a command are latched on the rising edge of WE# or CE# (whichever goes high first). Standard microprocessor write timings are used. Figures 18 and 19 illustrate WE# and CE# controlled write operations.

4 COMMAND DEFINITIONS

When the V_{CCW} voltage $\leq V_{CCWLK}$, read operations from the status register, identifier codes, or blocks are enabled. Placing $V_{CCWH1/2}$ on V_{CCW} enables successful block erase, full chip erase, word/byte write and lock-bit configuration operations.

Device operations are selected by writing specific commands into the CUI. Table 3 defines these commands.

	Tuere 2010 Due operations (DTTD::: + IH)								
Mode	Notes	RP#	CE#	OE#	WE#	Address	V _{CCW}	DQ ₀₋₁₅	RY/BY# ⁽³⁾
Read	8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	X	Х	D _{OUT}	X
Output Disable		V _{IH}	V _{IL}	V _{IH}	V _{IH}	X	Х	High Z	X
Standby		V _{IH}	V _{IH}	Х	Х	Х	Х	High Z	X
Reset	4	V _{IL}	Х	Х	X	X	Х	High Z	High Z
Read Identifier Codes	8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	See Figure 4, 5	Х	Note 5	High Z
Write	6,7,8	V _{IH}	V _{IL}	V_{IH}	V _{IL}	Х	Х	D _{IN}	X

Table 2.1. Bus Operations $(BYTE \#=V_{IH})^{(1,2)}$

Table 2.2. Bus Operations (BYTE#= V_{IL}) ^(1,2)									
Mode	Notes	RP#	CE#	OE#	WE#	Address	V _{CCW}	DQ ₀₋₇	RY/BY# ⁽³⁾
Read	8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	X	Х	D _{OUT}	Х
Output Disable		V _{IH}	V _{IL}	V _{IH}	V _{IH}	X	Х	High Z	X
Standby		V _{IH}	V _{IH}	Х	Х	X	Х	High Z	X
Reset	4	V _{IL}	X	Х	Х	X	Х	High Z	High Z
Read Identifier Codes	8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	See Figure 4, 5	Х	Note 5	High Z
Write	6,7,8	V _{IH}	V _{IL}	V _{IH}	V _{IL}	X	Х	D _{IN}	X

NOTES:

 Refer to DC Characteristics. When V_{CCW}≤V_{CCWLK}, memory contents can be read, but not altered.
 X can be V_{IL} or V_{IH} for control pins and addresses, and V_{CCWLK} or V_{CCWH1/2} for V_{CCW}. See DC Characteristics for V_{CCWLK} voltages.

3. RY/BY# is VOL when the WSM is executing internal block erase, full chip erase, word/byte write or lock-bit configuration algorithms. It is High Z during when the WSM is not busy, in block erase suspend mode (with word/byte write inactive), word/byte write suspend mode or reset mode.

4. RP# at GND±0.2V ensures the lowest power consumption.

5. See Section 4.2 for read identifier code data.

6. Command writes involving block erase, full chip erase, word/byte write or lock-bit configuration are reliably executed when $V_{CCW}=V_{CCWH1/2}$ and $V_{CC}=2.7V-3.6V$. 7. Refer to Table 3 for valid D_{IN} during a write operation.

8. Never hold OE# low and WE# low at the same timing.

(10)

Table 3. Command Definitions ^{(10)}								
	Bus Cycles		First Bus Cycle		Sec	Second Bus Cycle		
Command	Req'd.	Notes	Oper ⁽¹⁾	Addr ⁽²⁾	Data ⁽³⁾	Oper ⁽¹⁾	Addr ⁽²⁾	Data ⁽³⁾
Read Array/Reset	1		Write	Х	FFH			
Read Identifier Codes	≥2	4	Write	Х	90H	Read	IA	ID
Read Status Register	2		Write	Х	70H	Read	Х	SRD
Clear Status Register	1		Write	Х	50H			
Block Erase	2	5	Write	Х	20H	Write	BA	D0H
Full Chip Erase	2		Write	Х	30H	Write	Х	D0H
Word/Byte Write	2	5,6	Write	Х	40H or 10H	Write	WA	WD
Block Erase and Word/Byte Write Suspend	1	5	Write	Х	B0H			
Block Erase and Word/Byte Write Resume	1	5	Write	Х	D0H			
Set Block Lock-Bit	2	8	Write	Х	60H	Write	BA	01H
Clear Block Lock-Bits	2	7,8	Write	Х	60H	Write	Х	D0H
Set Permanent Lock-Bit	2	9	Write	Х	60H	Write	Х	F1H
OTP Program	2		Write	Х	C0H	Write	OA	OD

NOTES:

1. BUS operations are defined in Table 2.1 and Table 2.2.

2. X=Any valid address within the device. IA=Identifier Code Address: see Figure 4. BA=Address within the block being erased. WA=Address of memory location to be written. OA=Address of OTP block to be written: see Figure 5.

3. ID=Data read from identifier codes. SRD=Data read from status register. See Table 6 for a description of the status register bits. WD=Data to be written at location WA. Data is latched on the rising edge of WE# or CE# (whichever goes high first). OD=Data to be written at location OA. Data is latched on the rising edge of WE# or CE# (whichever goes high first).

- 4. Following the Read Identifier Codes command, read operations access manufacturer, device, block lock configuration and permanent lock configuration codes. See Section 4.2 for read identifier code data.
- 5. If WP# is V_{II} , boot blocks are locked without block lock-bits state. If WP# is V_{IH} , boot blocks are locked by block lockbits. The parameter and main blocks are locked by block lock-bits without WP# state.
- 6. Either 40H or 10H are recognized by the WSM as the word/byte write setup.
- 7. The clear block lock-bits operation simultaneously clears all block lock-bits.
- 8. If the permanent lock-bit is set, Set Block Lock-Bit and Clear Block Lock-Bits commands can not be done.
- 9. Once the permanent lock-bit is set, permanent lock-bit reset is unable.

10. Commands other than those shown above are reserved by SHARP for future device implementations and should not be used.

4.1 Read Array Command

Upon initial device power-up and after exit from reset mode, the device defaults to read array mode. This operation is also initiated by writing the Read Array command. The device remains enabled for reads until another command is written. Once the internal WSM has started a block erase, full chip erase, word/byte write or lock-bit configuration the device will not recognize the Read Array command until the WSM completes its operation unless the WSM is suspended via an Erase Suspend or Word/Byte Write Suspend command. The Read Array command functions independently of the V_{CCW} voltage and RP# can be V_{IH} .

4.2 Read Identifier Codes Command

The identifier code operation is initiated by writing the Read Identifier Codes command. Following the command write, read cycles from addresses shown in Figure 4 retrieve the manufacturer, device, block lock configuration and permanent lock configuration codes (see Table 4 for identifier code values). To terminate the operation, write another valid command. Like the Read Array command, the Read Identifier Codes command functions independently of the V_{CCW} voltage and RP# can be V_{IH}. Following the Read Identifier Codes command, the following information can be read:

 Table 4. Identifier Codes

Code	$\begin{array}{c} \text{Address}^{(2)} \\ \text{[A}_{20}\text{-}\text{A}_{0} \text{]} \end{array}$	$\begin{array}{c} Data^{(3)} \\ [DQ_7 \text{-} DQ_0] \end{array}$
Manufacture Code	00000H	B0H
Device Code	00001H	E3H
Block Lock Configuration	BA ⁽¹⁾ +2	
•Block is Unlocked		DQ ₀ =0
•Block is Locked		DQ ₀ =1
•Reserved for Future Use		DQ ₁₋₇
Permanent Lock Configuration	00003H	
•Device is Unlocked		DQ ₀ =0
•Device is Locked		DQ ₀ =1
•Reserved for Future Use		DQ ₁₋₇

NOTE:

- 1. BA selects the specific block lock configuration code to be read. See Figure 4 for the device identifier code memory map.
- 2. A_{-1} don't care in byte mode.
- 3. DQ_{15} - DQ_8 outputs 00H in word mode.

4.3 Read Status Register Command

The status register may be read to determine when a block erase, full chip erase, word/byte write or lock-bit configuration is complete and whether the operation completed successfully. It may be read at any time by writing the Read Status Register command. After writing this command, all subsequent read operations output data from the status register until another valid command is written. The status register contents are latched on the falling edge of OE# or CE#, whichever occurs. OE# or CE# must toggle to V_{IH} before further reads to update the status register latch. The Read Status Register command functions independently of the V_{CCW} voltage. RP# can be V_{IH} .

4.4 Clear Status Register Command

Status register bits SR.5, SR.4, SR.3 or SR.1 are set to "1"s by the WSM and can only be reset by the Clear Status Register command. These bits indicate various failure conditions (see Table 6). By allowing system software to reset these bits, several operations (such as cumulatively erasing multiple blocks or writing several words/bytes in sequence) may be performed. The status register may be polled to determine if an error occurred during the sequence.

To clear the status register, the Clear Status Register command (50H) is written. It functions independently of the applied V_{CCW} Voltage. RP# can be V_{IH} . This command is not functional during block erase or word/byte write suspend modes.

4.5 Block Erase Command

Erase is executed one block at a time and initiated by a two-cycle command. A block erase setup is first written, followed by an block erase confirm. This command sequence requires appropriate sequencing and an address within the block to be erased (erase changes all block data to FFFFH/FFH). Block preconditioning, erase, and verify are handled internally by the WSM (invisible to the system). After the two-cycle block erase sequence is written, the device automatically outputs status register data when read (see Figure 6). The CPU can detect block erase completion by analyzing the output data of the RY/BY# pin or status register bit SR.7.

When the block erase is complete, status register bit SR.5 should be checked. If a block erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read status register mode until a new command is issued.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Block Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable block erasure can only occur when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, block contents are protected against erasure. If block erase is attempted while $V_{CCW} \leq V_{CCWLK}$, SR.3 and SR.5 will be set to "1". Successful block erase requires for boot blocks that WP# is V_{IH} and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If block erase is attempted when the excepting above conditions, SR.1 and SR.5 will be set to "1".

4.6 Full Chip Erase Command

This command followed by a confirm command erases all of the unlocked blocks. A full chip erase setup (30H) is first written, followed by a full chip erase confirm (D0H). After a confirm command is written, device erases the all unlocked blocks block by block. This command sequence requires appropriate sequencing. Block preconditioning, erase and verify are handled internally by the WSM (invisible to the system). After the two-cycle full chip erase sequence is written, the device automatically outputs status register data when read (see Figure 7). The CPU can detect full chip erase completion by analyzing the output data of the RY/BY# pin or status register bit SR.7.

When the full chip erase is complete, status register bit SR.5 should be checked. If erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read

status register mode until a new command is issued. If error is detected on a block during full chip erase operation, WSM stops erasing. Full chip erase operation start from lower address block, finish the higher address block. Full chip erase can not be suspended.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Full Chip Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable full chip erasure can only occur when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, block contents are protected against erasure. If full chip erase is attempted while $V_{CCW} \leq V_{CCWLK}$, SR.3 and SR.5 will be set to "1". Successful full chip erase requires for boot blocks that WP# is V_{IH} and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If all blocks are locked, SR.1 and SR.5 will be set to "1".

4.7 Word/Byte Write Command

Word/Byte write is executed by a two-cycle command sequence. Word/Byte write setup (standard 40H or alternate 10H) is written, followed by a second write that specifies the address and data (latched on the rising edge of WE#). The WSM then takes over, controlling the word/byte write and write verify algorithms internally. After the word/byte write sequence is written, the device automatically outputs status register data when read (see Figure 8). The CPU can detect the completion of the word/byte write event by analyzing the RY/BY# pin or status register bit SR.7.

When word/byte write is complete, status register bit SR.4 should be checked. If word/byte write error is detected, the status register should be cleared. The internal WSM verify only detects errors for "1"s that do not successfully write to "0"s. The CUI remains in read status register mode until it receives another command.

Reliable word/byte writes can only occur when V_{CC}=2.7V-3.6V and V_{CCW}=V_{CCWH1/2}. In the absence of this high voltage, memory contents are protected against word/byte writes. If word/byte write is attempted while V_{CCW} \leq V_{CCWLK}, status register bits SR.3 and SR.4 will be set to "1". Successful word/byte write requires for boot blocks that WP# is V_{IH} and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If word/byte write is attempted when the excepting above conditions, SR.1 and SR.4 will be set to "1".

4.8 Block Erase Suspend Command

The Block Erase Suspend command allows block-erase interruption to read or word/byte write data in another block of memory. Once the block erase process starts, writing the Block Erase Suspend command requests that the WSM suspend the block erase sequence at a predetermined point in the algorithm. The device outputs status register data when read after the Block Erase Suspend command is written. Polling status register bits SR.7 and SR.6 can determine when the block erase operation has been suspended (both will be set to "1"). RY/BY# will also transition to High Z. Specification t_{WHRZ2} defines the block erase suspend latency.

When Block Erase Suspend command write to the CUI, if block erase was finished, the device places read array mode. Therefore, after Block Erase Suspend command write to the CUI, Read Status Register command (70H) has to write to CUI, then status register bit SR.6 should be checked for places the device in suspend mode.

At this point, a Read Array command can be written to read data from blocks other than that which is suspended. A Word/Byte Write command sequence can also be issued during erase suspend to program data in other blocks. Using the Word/Byte Write Suspend command (see Section 4.9), a word/byte write operation can also be suspended. During a word/byte write operation with block erase suspended, status register bit SR.7 will return to "0" and the RY/BY# output will transition to V_{OL} . However, SR.6 will remain "1" to indicate block erase suspend status.

The only other valid commands while block erase is suspended are Read Status Register and Block Erase Resume. After a Block Erase Resume command is written to the flash memory, the WSM will continue the block erase process. Status register bits SR.6 and SR.7 will automatically clear and RY/BY# will return to V_{OL} . After the Erase Resume command is written, the device automatically outputs status register data when read (see Figure 9). V_{CCW} must remain at $V_{CCWH1/2}$ (the same V_{CCW} level used for block erase) while block erase is suspended. RP# must also remain at V_{IL} or V_{IH} (the same WP# level used for block erase). Block erase cannot resume until word/byte write operations initiated during block erase suspend have completed.

If the time between writing the Block Erase Resume command and writing the Block Erase Suspend command is shorter than t_{ERES} and both commands are written repeatedly, a longer time is required than standard block erase until the completion of the operation.

4.9 Word/Byte Write Suspend Command

The Word/Byte Write Suspend command allows word/byte write interruption to read data in other flash memory locations. Once the word/byte write process starts, writing the Word/Byte Write Suspend command requests that the WSM suspend the Word/Byte write sequence at a predetermined point in the algorithm. The device continues to output status register data when read after the Word/Byte Write Suspend command is written. Polling status register bits SR.7 and SR.2 can determine when the word/byte write operation has been suspended (both will be set to "1"). RY/BY# will also transition to High Z. Specification t_{WHRZ1} defines the word/byte write suspend latency.

When Word/Byte Write Suspend command write to the CUI, if word/byte write was finished, the device places read array mode. Therefore, after Word/Byte Write Suspend command write to the CUI, Read Status Register command (70H) has to write to CUI, then status register bit SR.2 should be checked for places the device in suspend mode.

At this point, a Read Array command can be written to read data from locations other than that which is suspended. The only other valid commands while word/byte write is suspended are Read Status Register and Word/Byte Write Resume. After Word/Byte Write Resume command is written to the flash memory, the WSM will continue the word/byte write process. Status register bits SR.2 and SR.7 will automatically clear and RY/BY# will return to V_{OL} . After the Word/Byte Write Resume command is written, the device automatically outputs status register data when read (see Figure 10). V_{CCW} must remain at $V_{CCWH1/2}$ (the same V_{CCW} level used for word/byte write) while in word/byte write suspend mode. RP# must also remain at V_{IH} . WP# must also remain at V_{IL} or V_{IH} (the same WP# level used for word/byte write).

If the time between writing the Word/Byte Write Resume command and writing the Word/Byte Write Suspend command is short and both commands are written repeatedly, a longer time is required than standard word/byte write until the completion of the operation.

4.10 Set Block and Permanent Lock-Bit Commands

A flexible block locking and unlocking scheme is enabled via a combination of block lock-bits, a permanent lock-bit and WP# pin. The block lock-bits and WP# pin gates program and erase operations while the permanent lock-bit gates block-lock bit modification. With the permanent lock-bit not set, individual block lock-bits can be set using the Set Block Lock-Bit command. The Set Permanent Lock-Bit command, sets the permanent lock-bit. After the permanent lock-bit is set, block lock-bits and locked block contents cannot altered. See Table 5 for a summary of hardware and software write protection options.

Set block lock-bit and permanent lock-bit are executed by a two-cycle command sequence. The set block or permanent lock-bit setup along with appropriate block or device address is written followed by either the set block lock-bit confirm (and an address within the block to be locked) or the set permanent lock-bit confirm (and any device address). The WSM then controls the set lock-bit algorithm. After the sequence is written, the device automatically outputs status register data when read (see Figure 11). The CPU can detect the completion of the set lock-bit event by analyzing the RY/BY# pin output or status register bit SR.7.

When the set lock-bit operation is complete, status register bit SR.4 should be checked. If an error is detected, the status register should be cleared. The CUI will remain in read status register mode until a new command is issued.

This two-step sequence of set-up followed by execution ensures that lock-bits are not accidentally set. An invalid Set Block or Permanent Lock-Bit command will result in status register bits SR.4 and SR.5 being set to "1". Also, reliable operations occur only when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, lock-bit contents are protected against alteration.

A successful set block lock-bit operation requires that the permanent lock-bit be cleared. If it is attempted with the permanent lock-bit set, SR.1 and SR.4 will be set to "1" and the operation will fail.

4.11 Clear Block Lock-Bits Command

All set block lock-bits are cleared in parallel via the Clear Block Lock-Bits command. With the permanent lock-bit not set, block lock-bits can be cleared using only the Clear Block Lock-Bits command. If the permanent lock-bit is set, block lock-bits cannot cleared. See Table 5 for a summary of hardware and software write protection options.

Clear block lock-bits operation is executed by a two-cycle command sequence. A clear block lock-bits setup is first written. After the command is written, the device automatically outputs status register data when read (see Figure 12). The CPU can detect completion of the clear block lock-bits event by analyzing the RY/BY# Pin output or status register bit SR.7.

When the operation is complete, status register bit SR.5 should be checked. If a clear block lock-bit error is detected, the status register should be cleared. The CUI will remain in read status register mode until another command is issued.

This two-step sequence of set-up followed by execution ensures that block lock-bits are not accidentally cleared. An invalid Clear Block Lock-Bits command sequence will result in status register bits SR.4 and SR.5 being set to "1". Also, a reliable clear block lock-bits operation can only occur when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. If a clear block lock-bits operation is attempted while $V_{CCW}\leq V_{CCWLK}$, SR.3 and SR.5 will be set to "1". In the absence of this high voltage, the block lock-bits content are protected against alteration. A successful clear block lock-bits operation requires that the permanent lock-bit is not set. If it is attempted with the permanent lock-bit set, SR.1 and SR.5 will be set to "1" and the operation will fail.

If a clear block lock-bits operation is aborted due to V_{CCW} or V_{CC} transitioning out of valid range or RP# active transition, block lock-bit values are left in an undetermined state. A repeat of clear block lock-bits is required to initialize block lock-bit contents to known values. Once the permanent lock-bit is set, it cannot be cleared.

16

4.12 OTP Program Command

OTP program is executed by a two-cycle command sequence. OTP program command(C0H) is written, followed by a second write cycle that specifies the address and data (latched on the rising edge of WE#). The WSM then takes over, controlling the OTP program and program verify algorithms internally. After the OTP program command sequence is completed, the device automatically outputs status register data when read (see Figure 13). The CPU can detect the completion of the OTP program by analyzing the output data of the RY/BY# pin or status register bit SR.7.

When OTP program is completed, status register bit SR.4 should be checked. If OTP program error is detected, the status register should be cleared. The internal WSM verify only detects errors for "1"s that do not successfully program to "0"s. The CUI remains in read status register mode until it receives other commands.

Reliable OTP program can be executed only when V_{CC} =2.7V-3.6V and V_{CCW} = $V_{CCWH1/2}$. In the absence of this voltage, memory contents are protected against OTP

programs. If OTP program is attempted while $V_{CCW} \leq V_{CCWLK}$, status register bits SR.3 and SR.4 is set to "1". If OTP write is attempted when the OTP Lock-bit is set, SR.1 and SR.4 is set to "1".

4.13 Block Locking by the WP#

This Boot Block Flash memory architecture features two hardware-lockable boot blocks so that the kernel code for the system can be kept secure while other blocks are programmed or erased as necessary.

The lockable two boot blocks are locked when WP#=V_{IL}; any program or erase operation to a locked block will result in an error, which will be reflected in the status register. For top configuration, the top two boot blocks are lockable. For the bottom configuration, the bottom two boot blocks are lockable. If WP# is V_{IH} and block lockbit is not set, boot block can be programmed or erased normally (Unless V_{CCW} is below V_{CCWLK}). WP# is valid only two boot blocks, other blocks are not affected.

			100	10 5. WIII		
Operation	V _{CCW}	RP#	Permanent Lock-Bit	Block Lock-bit	WP#	Effect
Block Erase	≤V _{CCWLK}	Х	Х	X	X	All Blocks Locked.
or	>V _{CCWLK}	V _{IL}	Х	X	X	All Blocks Locked.
Word/Byte		V_{IH}	Х	0	V _{IL}	2 Boot Blocks Locked.
Write					V _{IH}	Block Erase and Word/Byte Write Enabled.
				1	V _{IL}	Block Erase and Word/Byte Write Disabled.
					V _{IH}	Block Erase and Word/Byte Write Disabled.
Full Chip	$\leq V_{CCWLK}$	Х	Х	Х	X	All Blocks Locked.
Erase	>V _{CCWLK}	V _{IL}	Х	Х	X	All Blocks Locked.
		V _{IH}	Х	X	V _{IL}	All Unlocked Blocks are Erased.
						2 Boot Blocks and Locked Blocks are NOT Erased.
					V _{IH}	All Unlocked Blocks are Erased,
						Locked Blocks are NOT Erased.
Set Block	$\leq V_{CCWLK}$	Х	Х	Х	X	Set Block Lock-Bit Disabled.
Lock-Bit	>V _{CCWLK}	V _{IL}	Х	Х	X	Set Block Lock-Bit Disabled.
		V_{IH}	0	Х	X	Set Block Lock-Bit Enabled.
			1	Х	X	Set Block Lock-Bit Disabled.
Clear Block	$\leq V_{CCWLK}$	Х	Х	X	X	Clear Block Lock-Bits Disabled.
Lock-Bits	>V _{CCWLK}	V _{IL}	Х	X	X	Clear Block Lock-Bits Disabled.
		V _{IH}	0	X	X	Clear Block Lock-Bits Enabled.
			1	Х	X	Clear Block Lock-Bits Disabled.
Set	≤V _{CCWLK}	Х	Х	Х	X	Set Permanent Lock-Bit Disabled.
Permanent	>V _{CCWLK}	V _{IL}	Х	Х	X	Set Permanent Lock-Bit Disabled.
Lock-Bit	· · ·	V _{IH}	Х	Х	X	Set Permanent Lock-Bit Enabled.

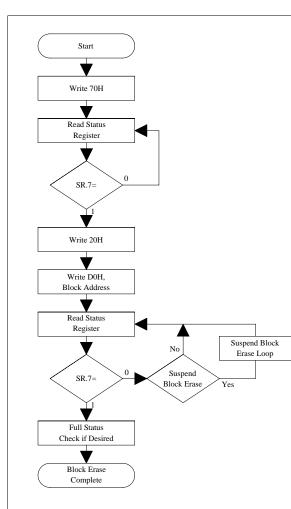
 Table 5. Write Protection Alternatives

SHARP

LHF32J08

WSMS	BESS	ECBLBS	WBWSLBS	VCCWS	WBWSS	DPS	R
7	6	5	4	3	2	1	0
					NOTES	:	
SR.7 = WRITH 1 = Ready 0 = Busy	E STATE MAC	CHINE STATUS	(WSMS)	erase, word/by	# or SR.7 to deter te write or lock- valid while SR.7=	bit configuratio	
1 = Block I	K ERASE SUS Erase Suspende Erase in Progres		S (BESS)				
STATU $1 = Error in$ $Lock-E$ $0 = Success$	JS (ECBLBS) n Block Erase, l Bits	R BLOCK LOCH Full Chip Erase o se, Full Chip Era	or Clear Block	erase or lock-b	nd SR.4 are "1"s bit configuration uence was entere	attempt, an imp	
 SR.4 = WORD/BYTE WRITE AND SET LOCK-BIT STATUS (WBWSLBS) 1 = Error in Word/Byte Write or Set Block/Permanent Lock-Bit 0 = Successful Word/Byte Write or Set Block/Permanent Lock-Bit 			level. The WS only after Bloc Lock-Bit Cont guaranteed to	provide a contin M interrogates a ck Erase, Full Ch figuration comma reports accurate	nd indicates the hip Erase, Word and sequences.	V _{CCW} level /Byte Write SR.3 is not	
SR.3 = V_{CCW} STATUS (VCCWS) 1 = V_{CCW} Low Detect, Operation Abort 0 = V_{CCW} OK					provide a contin		
SR.2 = WORD/BYTE WRITE SUSPEND STATUS (WBWSS) 1 = Word/Byte Write Suspended 0 = Word/Byte Write in Progress/Completed				the permanent Block Erase, F Configuration depending on	c-bit and WP# va lock-bit, block l Full Chip Erase, v command seque the attempted op	ock-bit and WP Word/Byte Writ nces. It informs eration, if the b	# only after te or Lock-Bi the system, lock lock-bit
 SR.1 = DEVICE PROTECT STATUS (DPS) 1 = Block Lock-Bit, Permanent Lock-Bit and/or WP# Lock Detected, Operation Abort 0 = Unlock 			the block lock writing the Re	t lock-bit is set an and permanent l ad Identifier Coc l block lock-bit s	ock configuration	on codes afte	
R.0 = RESEF	RVED FOR FU	TURE ENHAN	CEMENTS (R)		ed for future use he status register		nasked out





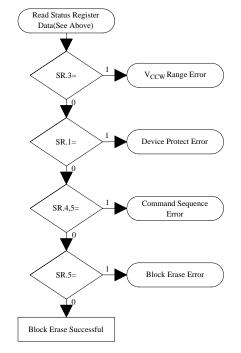
Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Erase Setup	Data=20H Addr=X
Write	Erase Confirm	Data=D0H Addr=Within Block to be Erased
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Danaat fan auhaaa	uant block arequires	L

Repeat for subsequent block erasures.

Full status check can be done after each block erase or after a sequence of block erasures.

Write FFH after the last operation to place device in read array mode.

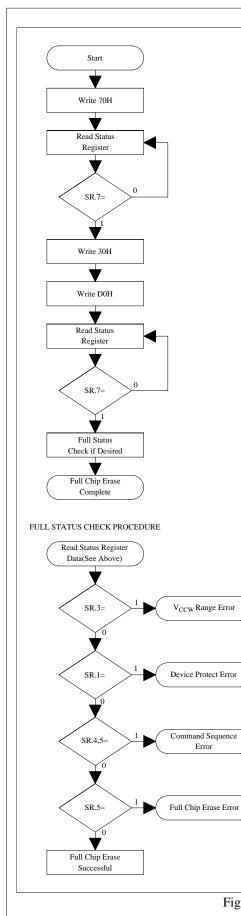
FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V _{CCW} Error Detect
Standby		Check SR.1 1=Device Protect Detect
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.5 1=Block Erase Error
where multiple	blocks are erased before full	y the Clear Status Register Command in cases I status is checked. 2fore attempting retry or other error recovery.

Figure 6. Automated Block Erase Flowchart

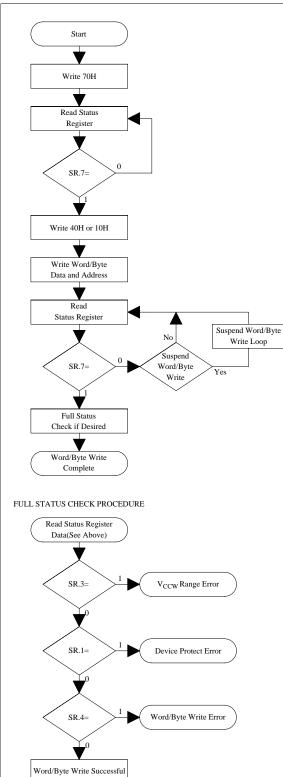




Command	Comments	
Read Status Register	Data=70H Addr=X	
	Status Register Data	
	Check SR.7 1=WSM Ready 0=WSM Busy	
Full Chip Erase Setup	Data=30H Addr=X	
Full Chip Erase Confirm	Data=D0H Addr=X	
	Status Register Data	
	Check SR.7 1=WSM Ready 0=WSM Busy	
	Read Status Register Full Chip Erase Setup Full Chip Erase	Read Status Register Data=70H Addr=X Status Register Data Status Register Data Check SR.7 1=WSM Ready 0=WSM Busy O=WSM Busy Full Chip Erase Setup Data=30H Addr=X Full Chip Erase Confirm Data=D0H Addr=X Status Register Data Status Register Data Check SR.7 1=WSM Ready Status Register Data

	Bus Operation	Command	Comments
Error	Standby		Check SR.3 1=V _{CCW} Error Detect
	Standby		Check SR.1 1=Device Protect Detect (All Blocks are locked)
t Error	Standby		Check SR.4,5 Both 1=Command Sequence Error
	Standby		Check SR.5 1=Full Chip Erase Error
juence	where multiple	blocks are erased before full	y the Clear Status Register Command in cases I status is checked. efore attempting retry or other error recovery.
e Error			
Figure 7. Automated	Full Chip E	rase Flowchart	





Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Setup Word/Byte Write	Data=40H or 10H Addr=X
Write	Word/Byte Write	Data=Data to Be Written Addr=Location to Be Written
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
SR full status che word/byte write	es.	ord/byte write, or after a sequence of ation to place device in read array mode.
Bus Operation	Command	Comments

V _{CCW} Range Error	\sum	
V CCW Kunge Enfor	\mathcal{I}	

Standby	Check SR.1 1=Device Protect Detect
Standby	Check SR.4 1=Data Write Error
where multiple locatio	only cleared by the Clear Status Register command in cases ns are written before full status is checked. the Status Register before attempting retry or other error recovery.

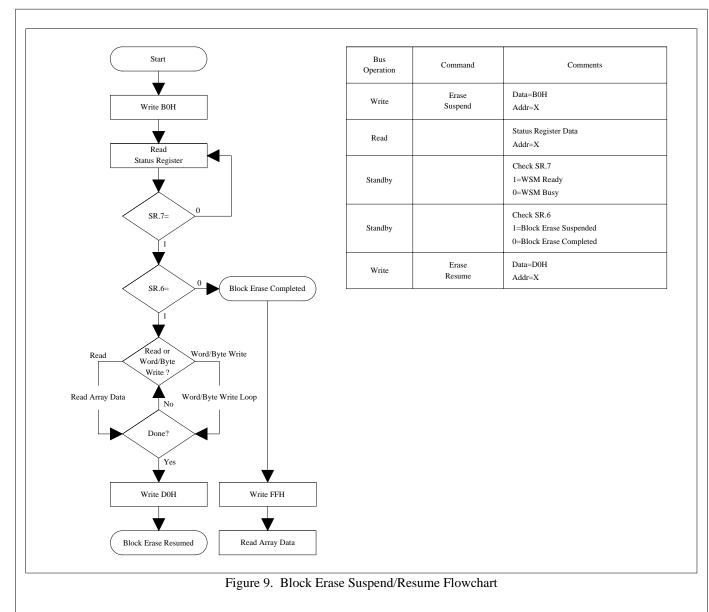
Check SR.3

 $1=V_{CCW}$ Error Detect

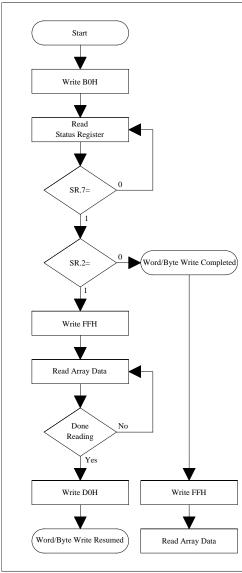
Figure 8. Automated Word/Byte Write Flowchart

Standby

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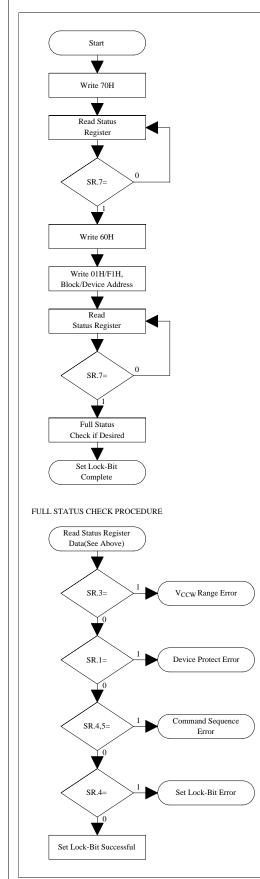




Command	Comments
Word/Byte Write Suspend	Data=B0H Addr=X
	Status Register Data Addr=X
	Check SR.7 1=WSM Ready 0=WSM Busy
	Check SR.2 1=Word/Byte Write Suspended 0=Word/Byte Write Completed
Read Array	Data=FFH Addr=X
	Read Array locations other than that being written.
Word/Byte Write Resume	Data=D0H Addr=X
	Word/Byte Write Suspend Read Array Word/Byte Write

Figure 10. Word/Byte Write Suspend/Resume Flowchart



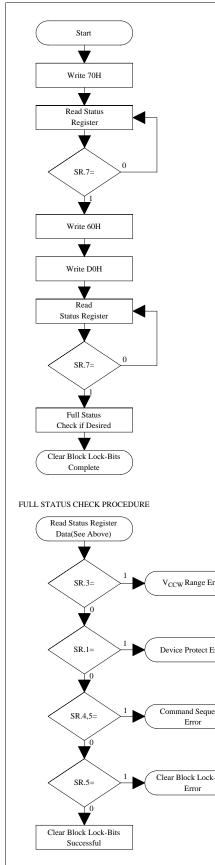


Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Set Block/Permanent Lock-Bit Setup	Data=60H Addr=X
Write	Set Block or Permanent Lock-Bit Confirm	Data=01H(Block), F1H(Permanent) Addr=Block Address(Block), Device Address(Permanent)
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Full status check lock-bit set ope	erations.	bit set operation or after a sequence of to place device in read array mode.

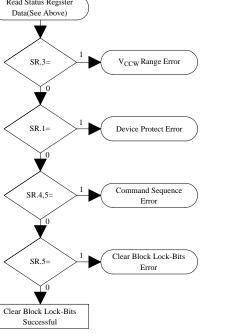
Bus Operation	Command	Comments
Standby		Check SR.3 1=V _{CCW} Error Detect
Standby		Check SR.1 1=Device Protect Detect Permanent Lock-Bit is Set (Set Block Lock-Bit Operation)
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.4 1=Set Lock-Bit Error
where multiple	lock-bits are set before full	y the Clear Status Register command in cases status is checked. efore attempting retry or other error recovery.

Figure 11. Set Block and Permanent Lock-Bit Flowchart





Bus Operation	Command	Comments			
Write	Read Status Register	Data=70H Addr=X			
Read		Status Register Data			
Standby		Check SR.7 1=WSM Ready 0=WSM Busy			
Write	Clear Block Lock-Bits Setup	Data=60H Addr=X			
Write	Clear Block Lock-Bits Confirm	Data=D0H Addr=X			
Read		Status Register Data			
Standby		Check SR.7 1=WSM Ready 0=WSM Busy			

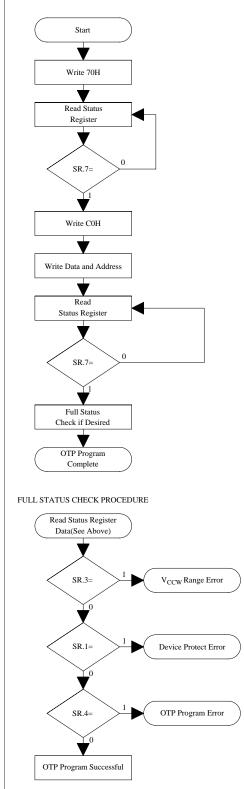


Bus Operation	Command	Comments
Standby		Check SR.3 1=V _{CCW} Error Detect
Standby		Check SR.1 1=Device Protect Detect Permanent Lock-Bit is Set
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.5 1=Clear Block Lock-Bits Error
		y the Clear Status Register command. efore attempting retry or other error recovery.



25

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Bus Operation	Command	Comments				
Write	Read Status Register	Data=70H Addr=X				
Read		Status Register Data				
Standby		Check SR.7 1=WSM Ready 0=WSM Busy				
Write	Setup OTP Program	Data=C0H Addr=X				
Write	OTP Program	Data=Data to Be Written Addr=Location to Be Written				
Read		Status Register Data				
Standby		Check SR.7 1=WSM Ready 0=WSM Busy				
SR full status che OTP programs		TP program, or after a sequence of on to place device in read array mode.				
Bus Operation Command Comments						
		Check SR 3				

SR.3= V _{CCW} Range Error
SR.1= Device Protect Error
\mathbf{X}^{0}
SR.4= 1 OTP Program Error
Program Successful
PTO9TAILI SUCCESSIUL

Bus Operation	Command	Comments					
Standby		Check SR.3 1=V _{CCW} Error Detect					
Standby		Check SR.1 1=Device Protect Detect					
Standby		Check SR.4 1=Data Write Error					
where multiple	locations are written before	Clear Status Register command in cases full status is checked. efore attempting retry or other error recovery.					

Figure 13. Automated OTP Program Flowchart

5 DESIGN CONSIDERATIONS

5.1 Three-Line Output Control

The device will often be used in large memory arrays. SHARP provides three control inputs to accommodate multiple memory connections. Three-line control provides for:

- a. Lowest possible memory power dissipation.
- b. Complete assurance that data bus contention will not occur.

To use these control inputs efficiently, an address decoder should enable CE# while OE# should be connected to all memory devices and the system's READ# control line. This assures that only selected memory devices have active outputs while deselected memory devices are in standby mode. RP# should be connected to the system POWERGOOD signal to prevent unintended writes during system power transitions. POWERGOOD should also toggle during system reset.

5.2 RY/BY# and WSM Polling

RY/BY# is an open drain output that should be connected to V_{CC} by a pull up resistor to provides a hardware method of detecting block erase, full chip erase, word/byte write and lock-bit configuration completion. It transitions low after block erase, full chip erase, word/byte write or lock-bit configuration commands and returns to V_{OH} (while RY/BY# is pull up) when the WSM has finished executing the internal algorithm.

RY/BY# can be connected to an interrupt input of the system CPU or controller. It is active at all times. RY/BY# is also High Z when the device is in block erase suspend (with word/byte write inactive), word/byte write suspend or reset modes.

5.3 Power Supply Decoupling

Flash memory power switching characteristics require careful device decoupling. System designers are interested in three supply current issues; standby current levels, active current levels and transient peaks produced by falling and rising edges of CE# and OE#. Transient current magnitudes depend on the device outputs' capacitive and inductive loading. Two-line control and proper decoupling capacitor selection will suppress transient voltage peaks. Each device should have a 0.1µF ceramic capacitor connected between its V_{CC} and GND and between its V_{CCW} and GND. These high-frequency, low inductance capacitors should be placed as close as possible to package leads. Additionally, for every eight devices, a 4.7µF electrolytic capacitor should be placed at the array's power supply connection between V_{CC} and GND. The bulk capacitor will overcome voltage slumps caused by PC board trace inductance.

5.4 V_{CCW} Trace on Printed Circuit Boards

Updating flash memories that reside in the target system requires that the printed circuit board designer pay attention to the V_{CCW} Power supply trace. The V_{CCW} pin supplies the memory cell current for word/byte writing and block erasing. Use similar trace widths and layout considerations given to the V_{CC} power bus. Adequate V_{CCW} supply traces and decoupling will decrease V_{CCW} voltage spikes and overshoots.

5.5 V_{CC}, V_{CCW}, RP# Transitions

Block erase, full chip erase, word/byte write and lock-bit configuration are not guaranteed if V_{CCW} falls outside of a valid $V_{CCWH1/2}$ range, V_{CC} falls outside of a valid 2.7V-3.6V range, or RP# $\neq V_{IH}$. If V_{CCW} error is detected, status register bit SR.3 is set to "1" along with SR.4 or SR.5, depending on the attempted operation. If RP# transitions to V_{IL} during block erase, full chip erase, word/byte write or lock-bit configuration, RY/BY# will remain low until the reset operation is complete. Then, the operation will abort and the device will enter reset mode. The aborted operation may leave data partially altered. Therefore, the command sequence must be repeated after normal operation is restored. Device power-off or RP# transitions to V_{IL} clear the status register.

The CUI latches commands issued by system software and is not altered by V_{CCW} or CE# transitions or WSM actions. Its state is read array mode upon power-up, after exit from reset mode or after V_{CC} transitions below V_{LKO} .

5.6 Power-Up/Down Protection

The device is designed to offer protection against accidental block erase, full chip erase, word/byte write or lock-bit configuration during power transitions. Upon power-up, the device is indifferent as to which power supply (V_{CCW} or V_{CC}) powers-up first. Internal circuitry resets the CUI to read array mode at power-up.

A system designer must guard against spurious writes for V_{CC} voltages above V_{LKO} when V_{CCW} is active. Since both WE# and CE# must be low for a command write, driving either to V_{IH} will inhibit writes. The CUI's two-step command sequence architecture provides added level of protection against data alteration.

In-system block lock and unlock capability prevents inadvertent data alteration. The device is disabled while $RP\#=V_{IL}$ regardless of its control inputs state.

5.7 Power Dissipation

When designing portable systems, designers must consider battery power consumption not only during device operation, but also for data retention during system idle time. Flash memory's nonvolatility increases usable battery life because data is retained when system power is removed.

5.8 Data Protection Method

Noises having a level exceeding the limit specified in the specification may be generated under specific operating conditions on some systems. Such noises, when induced onto WE# signal or power supply, may be interpreted as false commands, causing undesired memory updating. To protect the data stored in the flash memory against unwanted overwriting, systems operating with the flash memory should have the following write protect designs, as appropriate:

1) Protecting data in specific block

When a lock bit is set, the corresponding block (includes the 2 boot blocks) is protected against overwriting. By setting a WP# to low, only the 2 boot blocks can be protected against overwriting. By using this feature, the flash memory space can be divided into the program section (locked section) and data section (unlocked section). The permanent lock bit can be used to prevent false block bit setting. For further information on setting/resetting lock-bit, refer to the specification. (See chapter 4.10 and 4.11.)

2) Data protection through V_{CCW}

When the level of V_{CCW} is lower than V_{CCWLK} (lockout voltage), write operation on the flash memory is disabled. All blocks are locked and the data in the blocks are completely write protected. For the lockout voltage, refer to the specification. (See chapter 6.2.3.)

3) Data protection through RP#

When the RP# is kept low during read mode, the flash memory will be reset mode, then write protecting all blocks. When the RP# is kept low during power up and power down sequence such as voltage transition, write operation on the flash memory is disabled, write protecting all blocks. For the details of RP# control, refer to the specification. (See chapter 5.6 and 6.2.7.)

6 ELECTRICAL SPECIFICATIONS

6.1 Absolute Maximum Ratings*

*WARNING: Stressing the device beyond the "Absolute Maximum Ratings" may cause permanent damage. These are stress ratings only. Operation beyond the "Operating Conditions" is not recommended and extended exposure beyond the "Operating Conditions" may affect device reliability.

NOTES:

- 1. Operating temperature is for extended temperature product defined by this specification.
- 2. All specified voltages are with respect to GND. Minimum DC voltage is -0.5V on input/output pins and -0.2V on V_{CC} and V_{CCW} pins. During transitions, this level may undershoot to -2.0V for periods <20ns. Maximum DC voltage on input/output pins are V_{CC} +0.5V which, during transitions, may overshoot to V_{CC} +2.0V for periods <20ns.
- 3. Maximum DC voltage on V_{CCW} may overshoot to +13.0V for periods <20ns. Applying $12V\pm0.3V$ to V_{CCW} during erase/write can only be done for a maximum of 1000 cycles on each block. V_{CCW} may be connected to $12V\pm0.3V$ for a total of 80 hours maximum.
- 4. Output shorted for no more than one second. No more than one output shorted at a time.

6.2 Operating Conditions

Temperature and V _{CC} Operating Conditions							
Symbol	Parameter	Min.	Max.	Unit	Test Condition		
T _A	Operating Temperature	-40	+85	°C	Ambient Temperature		
V _{CC}	V _{CC} Supply Voltage (2.7V-3.6V)	2.7	3.6	V			

6.2.1 CAPACITANCE⁽¹⁾

T _A =+25°C, f=1MHz						
Symbol	Parameter	Тур.	Max.	Unit	Condition	
C _{IN}	Input Capacitance	7	10	pF	V _{IN} =0.0V	
C _{OUT}	Output Capacitance	9	12	pF	V _{OUT} =0.0V	
		J	í	1	001	

NOTE:

1. Sampled, not 100% tested.

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6.2.2 AC INPUT/OUTPUT TEST CONDITIONS

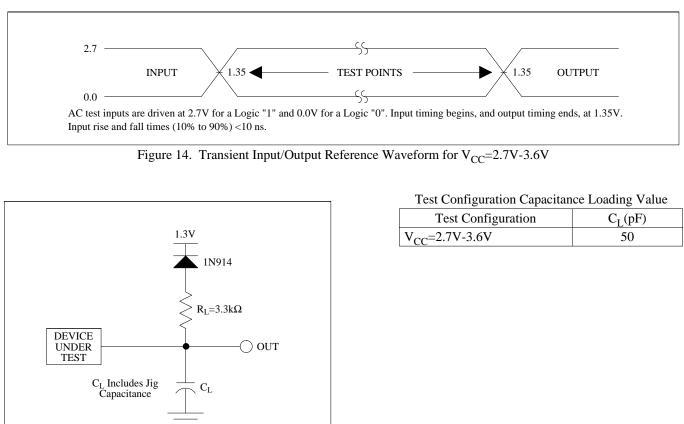


Figure 15. Transient Equivalent Testing Load Circuit

6.2.3 DC CHARACTERISTICS

DC Characteristics						
			V _{CC} =2.	7V-3.6V		Test
Sym.	Parameter	Notes	Тур.	Max.	Unit	Conditions
I _{LI}	Input Load Current	1		±0.5	μΑ	V _{CC} =V _{CC} Max. V _{IN} =V _{CC} or GND
I _{LO}	Output Leakage Current	1		±0.5	μΑ	V _{CC} =V _{CC} Max. V _{OUT} =V _{CC} or GND
I _{CCS}	V _{CC} Standby Current	1,3	4	20	μΑ	V _{CC} =V _{CC} Max. CE#=RP#=V _{CC} ±0.2V
I _{CCAS}	V _{CC} Auto Power-Save Current	1,5	4	20	μΑ	V _{CC} =V _{CC} Max. CE#=GND±0.2V
I _{CCD}	V _{CC} Reset Power-Down Current	1	4	20	μΑ	RP#=GND±0.2V I _{OUT} (RY/BY#)=0mA
I _{CCR}	V _{CC} Read Current	1	15	30	mA	V _{CC} =V _{CC} Max., CE#=GND f=5MHz, I _{OUT} =0mA
I _{CCW}	V _{CC} Word/Byte Write or Set Lock-	1,6	5	17	mA	V _{CCW} =2.7V-3.6V
	Bit Current		5	12	mA	V _{CCW} =11.7V-12.3V
I _{CCE}	V _{CC} Block Erase, Full Chip Erase or	1,6	4	17	mA	V _{CCW} =2.7V-3.6V
	Clear Block Lock-Bits Current		4	12	mA	V _{CCW} =11.7V-12.3V
I _{CCWS} I _{CCES}	V _{CC} Word/Byte Write or Block Erase Suspend Current	1,2	1	6	mA	CE#=V _{IH}
I _{CCWS}	V _{CCW} Standby or Read Current	1	±2	±15	μΑ	V _{CCW} ≤V _{CC}
I _{CCWR}			10	200	μΑ	V _{CCW} >V _{CC}
I _{CCWAS}	V _{CCW} Auto Power-Save Current	1,5	0.1	5	μΑ	V _{CC} =V _{CC} Max. CE#=GND±0.2V
I _{CCWD}	V _{CCW} Reset Power-Down Current	1	0.1	5	μΑ	RP#=GND±0.2V
I _{CCWW}	V _{CCW} Word/Byte Write or Set Lock-	1,6	12	40	mA	V _{CCW} =2.7V-3.6V
	Bit Current			30	mA	V _{CCW} =11.7V-12.3V
I _{CCWE}	V _{CCW} Block Erase, Full Chip Erase	1,6	8	25	mA	V _{CCW} =2.7V-3.6V
	or Clear Block Lock-Bits Current			20	mA	V _{CCW} =11.7V-12.3V
I _{CCWWS} I _{CCWES}	V _{CCW} Word/Byte Write or Block Erase Suspend Current	1	10	200	μΑ	V _{CCW} =V _{CCWH1/2}

DC Characteristics (Continued)						
			V _{CC} =2.7V-3.6V			
Sym.	Parameter	Notes	Min.	Max.	Unit	Test Conditions
V _{IL}	Input Low Voltage	6	-0.5	0.4	V	
V _{IH}	Input High Voltage	6	V _{CC} -0.4	V _{CC} +0.5	v	
V _{OL}	Output Low Voltage	3,6		0.4	v	V _{CC} =V _{CC} Min. I _{OL} =2.0mA
V _{OH}	Output High Voltage	6	V _{CC} -0.4		v	$V_{CC}=V_{CC}$ Min. $I_{OL}=2.0$ mA $V_{CC}=V_{CC}$ Min. $I_{OH}=-100\mu$ A
V _{CCWLK}	V _{CCW} Lockout during Normal Operations	4,6		1.0	v	
V _{CCWH1}	V _{CCW} during Block Erase, Full Chip Erase, Word/Byte Write or Lock-Bit Configuration Operations		2.7	3.6	v	
V _{CCWH2}	V _{CCW} during Block Erase, Full Chip Erase, Word/Byte Write or Lock-Bit Configuration Operations	7	11.7	12.3	v	
V _{LKO}	V _{CC} Lockout Voltage		2.0		V	

NOTES:

1. All currents are in RMS unless otherwise noted. Typical values at nominal V_{CC} voltage and T_A =+25°C.

2. I_{CCWS} and I_{CCES} are specified with the device de-selected. If read or word/byte written while in erase suspend mode, the device's current draw is the sum of I_{CCWS} or I_{CCES} and I_{CCR} or I_{CCW}, respectively.

3. Includes RY/BY#.

Block erases, full chip erase, word/byte writes and lock-bit configurations are inhibited when V_{CCW}≤V_{CCWLK}, and not guaranteed in the range between V_{CCWLK}(max.) and V_{CCWH1}(min.), between V_{CCWH1}(max.) and V_{CCWH2}(min.) and above V_{CCWH2}(max.).

5. The Automatic Power Savings (APS) feature is placed automatically power save mode that addresses not switching more than 300ns while read mode.

6. Sampled, not 100% tested.

7. Applying 12V±0.3V to V_{CCW} during erase/write can only be done for a maximum of 1000 cycles on each block. V_{CCW} may be connected to 12V±0.3V for a total of 80 hours maximum.

32

6.2.4 AC CHARACTERISTICS - READ-ONLY OPERATIONS $^{(1)}$

V_{CC} =2.7V-3.6V, T_{A} =-40°C to +85°C						
Sym.	Parameter	Notes	Min.	Max.	Unit	
t _{AVAV}	Read Cycle Time		90		ns	
t _{AVQV}	Address to Output Delay			90	ns	
t _{ELQV}	CE# to Output Delay	2		90	ns	
t _{PHOV}	RP# High to Output Delay			600	ns	
t _{GLQV}	OE# to Output Delay	2		40	ns	
t _{ELQX}	CE# to Output in Low Z	3	0		ns	
t _{EHQZ}	CE# High to Output in High Z	3		40	ns	
t _{GLOX}	OE# to Output in Low Z	3	0		ns	
t _{GHQZ}	OE# High to Output in High Z	3		15	ns	
t _{OH}	Output Hold from Address, CE# or OE# Change, Whichever Occurs First	3	0		ns	
t _{FVOV}	BYTE# to Output Delay	3		90	ns	
t _{FLQZ}	BYTE# Low to Output in High Z	3		25	ns	
t _{ELFV}	CE# to BYTE# High or Low	3,4		5	ns	

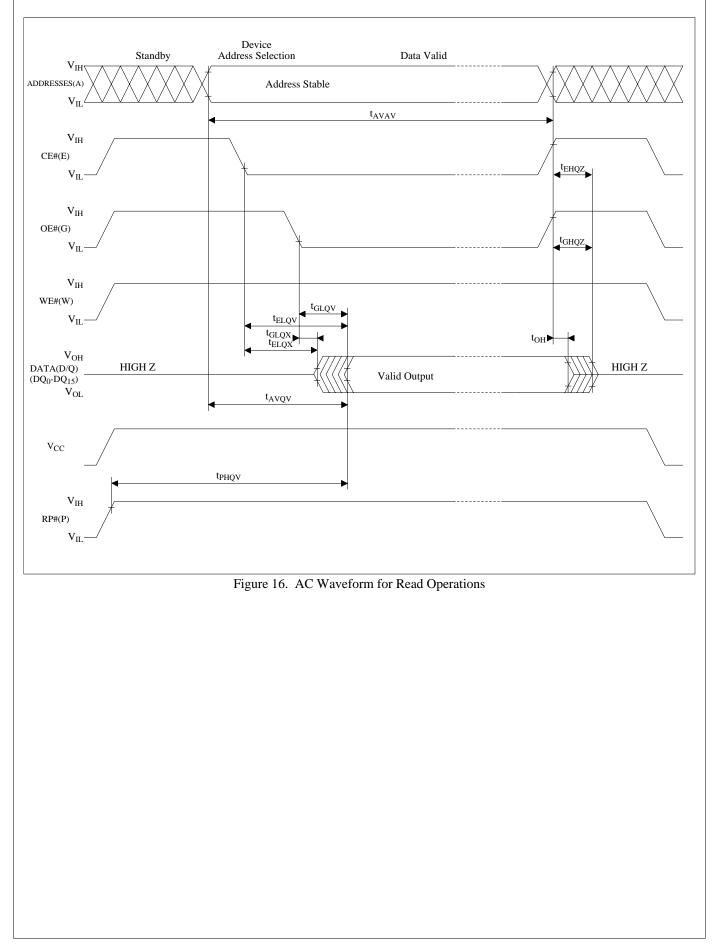
NOTES:

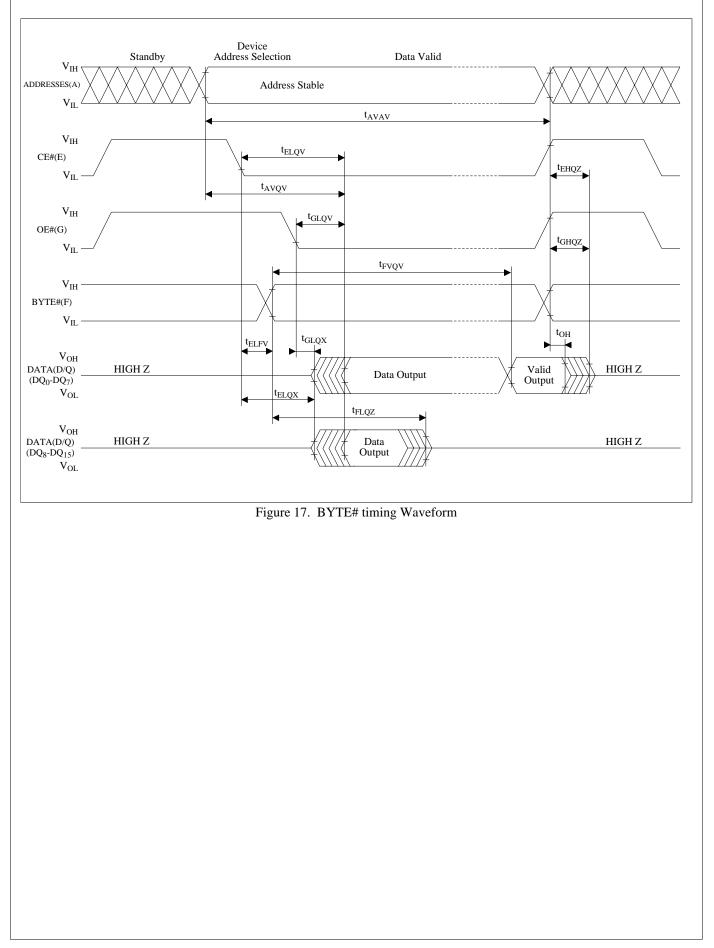
1. See AC Input/Output Reference Waveform for maximum allowable input slew rate.

2. OE# may be delayed up to t_{ELOV} - t_{GLOV} after the falling edge of CE# without impact on t_{ELOV} .

3. Sampled, not 100% tested.

4. If BYTE# transfer during reading cycle, exist the regulations separately.





36

6.2.5 AC CHARACTERISTICS - WRITE OPERATIONS ⁽¹⁾	6.2.5	AC CHARACTERISTICS -	WRITE OPERATIONS ⁽¹⁾
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Sym.	Parameter	Notes	Min.	Max.	Unit
t _{AVAV}	Write Cycle Time		90		ns
t _{PHWL}	RP# High Recovery to WE# Going Low	2	1		μs
t _{ELWL}	CE# Setup to WE# Going Low		10		ns
t _{WLWH}	WE# Pulse Width		50		ns
t _{SHWH}	WP#V _{IH} Setup to WE# Going High	2	100		ns
t _{VPWH}	V _{CCW} Setup to WE# Going High	2	100		ns
t _{AVWH}	Address Setup to WE# Going High	3	50		ns
t _{DVWH}	Data Setup to WE# Going High	3	50		ns
t _{WHDX}	Data Hold from WE# High		0		ns
t _{WHAX}	Address Hold from WE# High		0		ns
t _{WHEH}	CE# Hold from WE# High		10		ns
t _{WHWL}	WE# Pulse Width High		30		ns
t _{WHRL}	WE# High to RY/BY# Going Low or SR.7 Going "0"			100	ns
t _{WHGL}	Write Recovery before Read		0		ns
t _{QVVL}	V _{CCW} Hold from Valid SRD, RY/BY# High Z	2,4	0		ns
t _{OVSL}	WP# V _{IH} Hold from Valid SRD, RY/BY# High Z	2,4	0		ns
t _{FVWH}	BYTE# Setup to WE# Going High	5	50		ns
t _{WHFV}	BYTE# Hold from WE# High	5	90		ns

711 2 (11 T 1000

NOTES:

1. Read timing characteristics during block erase, full chip erase, word/byte write and lock-bit configuration operations are the same as during read-only operations. Refer to AC Characteristics for read-only operations.

2. Sampled, not 100% tested.

Balipica, not room condition.
 Refer to Table 4 for valid A_{IN} and D_{IN} for block erase, full chip erase, word/byte write or lock-bit configuration.
 V_{CCW} should be held at V_{CCWH1/2} until determination of block erase, full chip erase, word/byte write or lock-bit configuration success (SR.1/3/4/5=0).

5. If BYTE# switch during reading cycle, exist the regulations separately.

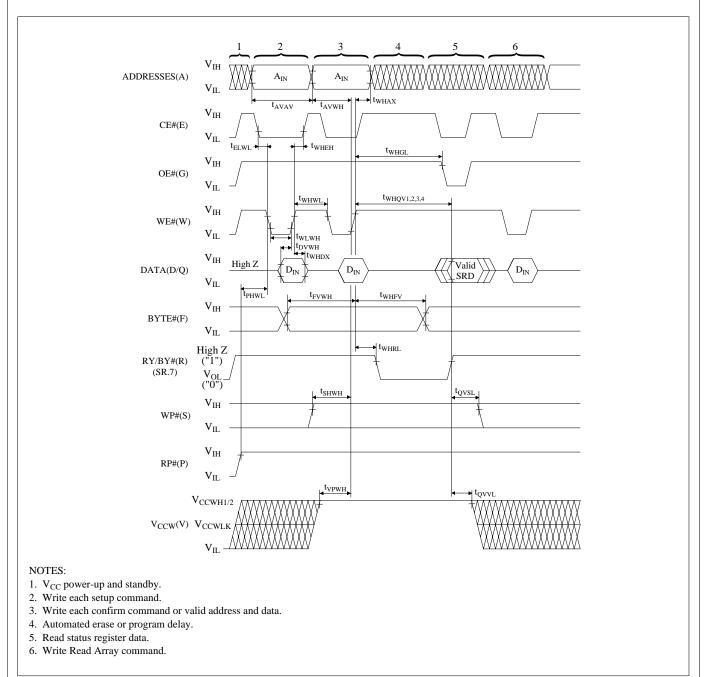


Figure 18. AC Waveform for WE#-Controlled Write Operations

6.2.6 ALTERNATIVE CE#-CONTROLLED WRITES⁽¹⁾

Sym.	Parameter	Notes	Min.	Max.	Unit
t _{AVAV}	Write Cycle Time		90		ns
t _{PHEL}	RP# High Recovery to CE# Going Low	2	1		μs
t _{WLEL}	WE# Setup to CE# Going Low		0		ns
t _{ELEH}	CE# Pulse Width		65		ns
t _{SHEH}	WP#V _{IH} Setup to CE# Going High	2	100		ns
t _{VPEH}	V _{CCW} Setup to CE# Going High	2	100		ns
t _{AVEH}	Address Setup to CE# Going High	3	50		ns
t _{DVEH}	Data Setup to CE# Going High	3	50		ns
t _{EHDX}	Data Hold from CE# High		0		ns
t _{EHAX}	Address Hold from CE# High		0		ns
t _{EHWH}	E# Hold from CE# High 0			ns	
t _{EHEL}	CE# Pulse Width High		25		ns
t _{EHRL}	CE# High to RY/BY# Going Low or SR.7 Going "0"			100	ns
t _{EHGL}	Write Recovery before Read		0		ns
t _{QVVL}	V _{CCW} Hold from Valid SRD, RY/BY# High Z	2,4	0		ns
t _{QVSL}	WP# V _{IH} Hold from Valid SRD, RY/BY# High Z	2,4	0		ns
t _{FVEH}	BYTE# Setup to CE# Going High	5	50		ns
t _{EHFV}	BYTE# Hold from CE# High	5	90		ns

711 2 (11 T 1000

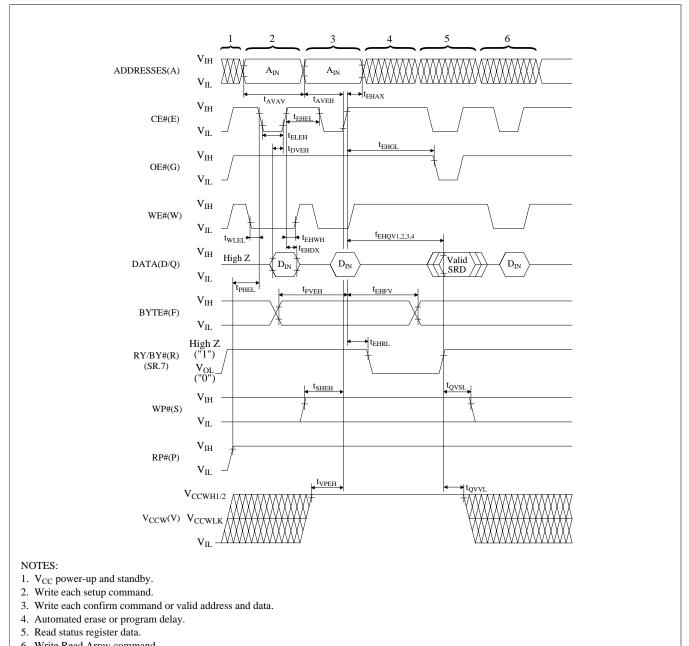
NOTES:

1. In systems where CE# defines the write pulse width (within a longer WE# timing waveform), all setup, hold, and inactive WE# times should be measured relative to the CE# waveform.

2. Sampled, not 100% tested.

Refer to Table 4 for valid A_{IN} and D_{IN} for block erase, full chip erase, word/byte write or lock-bit configuration.
 V_{CCW} should be held at V_{CCWH1/2} until determination of block erase, full chip erase, word/byte write or lock-bit configuration success (SR.1/3/4/5=0).

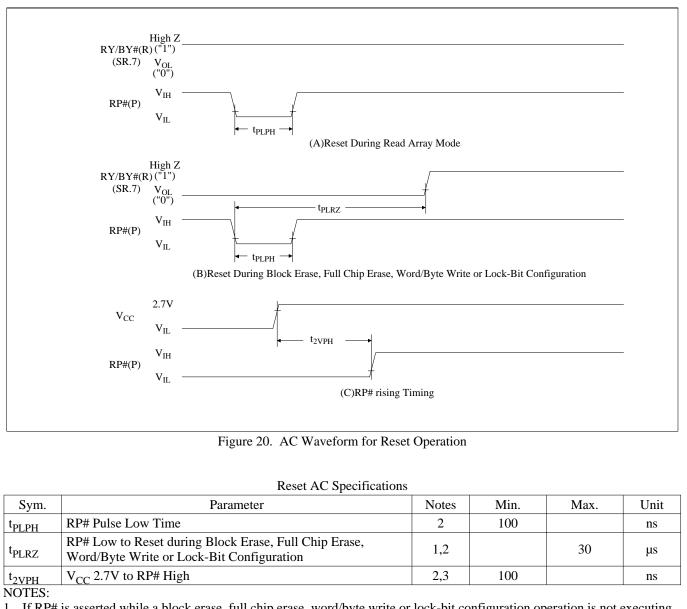
5. If BYTE# switch during reading cycle, exist the regulations separately.







6.2.7 RESET OPERATIONS



1. If RP# is asserted while a block erase, full chip erase, word/byte write or lock-bit configuration operation is not executing, the reset will complete within 100ns.

A reset time, t_{PHQV}, is required from the later of RY/BY#(SR.7) going High Z("1") or RP# going high until outputs are valid. Refer to AC Characteristics - Read-Only Operations for t_{PHOV}.

3. When the device power-up, holding RP# low minimum 100ns is required after V_{CC} has been in predefined range and also has been in stable there.

6.2.8	BLOCK ERASE, FULL CHIP ERASE, WORD/BYTE WRITE AND LOCK-BIT
	CONFIGURATION PERFORMANCE ⁽³⁾

$V_{CC}=2.7V-3.6V, T_{A}=-40^{\circ}C \text{ to } +85^{\circ}C$										
				V _{CCW} =2.7V-3.6V V _{CCW} =11.7V-12.3V			12.3V			
Sym.	Parame	eter	Notes	Min.	Typ. ⁽¹⁾	Max.	Min.	Typ. ⁽¹⁾	Max.	Unit
t _{WHQV1}	Word Write Time	32K word Block	2		33	200		20		μs
t _{EHQV1}		4K word Block	2		36	200		27		μs
	Byte Write Time	64K byte Block	2		31	200		19		μs
		8K byte Block	2		32	200		26		μs
	Block Write Time	32K word Block	2		1.1	4		0.66		s
	(In word mode)	4K word Block	2		0.15	0.5		0.12		s
	Block Write Time	64K byte Block	2		2.2	7		1.4		s
	(In byte mode)	8K byte Block	2		0.3	1		0.25		s
t _{WHQV2} t _{EHQV2}	Block Erase Time	32K word Block 64K byte Block	2		1.2	6		0.9		s
		4K word Block 8K byte Block	2		0.6	5		0.5		s
	Full Chip Erase Time		2		84	420		64		s
t _{WHQV3} t _{EHQV3}	Set Lock-Bit Time		2		56	200		42		μs
t _{WHQV4} t _{EHQV4}	Clear Block Lock-Bits	Time	2		1	5		0.69		s
t _{WHRZ1} t _{EHRZ1}	Word/Byte Write Suspe Read	nd Latency Time to	4		6	15		6	15	μs
t _{WHRZ2} t _{EHRZ2}	Block Erase Suspend La Read	atency Time to	4		16	30		16	30	μs
t _{ERES}	Latency Time from Bloc Command to Block Eras Command		5	600			600			μs

NOTES:

1. Typical values measured at $T_A = +25^{\circ}C$ and $V_{CC} = 3.0V$, $V_{CCW} = 3.0V$ or 12.0V. Assumes corresponding lock-bits are not set. Subject to change based on device characterization.

2. Excludes system-level overhead.

3. Sampled but not 100% tested.

4. A latency time is required from issuing suspend command(WE# or CE# going high) until RY/BY# going High Z or SR.7 going "1".

5. If the time between writing the Block Erase Resume command and writing the Block Erase Suspend command is shorter than t_{ERES} and both commands are written repeatedly, a longer time is required than standard block erase until the completion of the operation.



7 Package and packing specification

1.Storage Conditions.

- 1-1. Storage conditions required before opening the dry packing.
 - Normal temperature $:5{\sim}40^{\circ}$
 - Normal humidity : 80% R.H. max.

1-2. Storage conditions required after opening the dry packing.

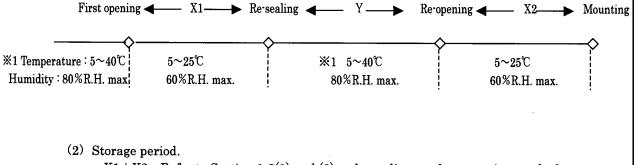
In order to prevent moisture absorption after opening, ensure the following storage conditions apply:

- (1) Storage conditions for one-time soldering. (Convection reflow, IR/Convection reflow, or Manual soldering.)
 - Temperature $:5 \sim 25$ °C
 - Humidity : 60% R.H. max.
 - Period : 72 hours max. after opening.
- (2) Storage conditions for two-time soldering. (Convection reflow, IR/Convection reflow.)
 - a. Storage conditions following opening and prior to performing the 1st reflow.
 - Temperature $: 5 \sim 25$ °C.
 - Humidity : 60% R.H. max.
 - Period : 72 hours max. after opening.
 - b. Storage conditions following completion of the 1st reflow and prior to performing the 2nd reflow.
 - Temperature $: 5 \sim 25^{\circ} C_{\circ}$
 - Humidity : 60% R.H. max.
 - Period : 72 hours max. after completion of the 1st reflow.

1-3. Temporary storage after opening.

To re-store the devices before soldering, do so only once and use a dry box or place desiccant (with a blue humidity indicator) with the devices and perform dry packing again using heat-sealing.

- The storage period, temperature and humidity must be as follows :
- (1) Storage temperature and humidity.
 - [∞]1 : External atmosphere temperature and humidity of the dry packing.



• X1 + X2: Refer to Section 1-2(1) and (2)a , depending on the mounting method.

• Y : Two weeks max.

- 2. Baking Condition.
 - (1) Situations requiring baking before mounting.
 - Storage conditions exceed the limits specified in Section 1-2 or 1-3.
 - Humidity indicator in the desiccant was already pink when opened.
 - (Also for re-opening.)
 - (2) Recommended baking conditions.
 - Baking temperature and period :
 - 120° for $16 \sim 24$ hours.
 - The above baking conditions apply since the trays are heat-resistant.
 - (3) Storage after baking.
 - After baking, store the devices in the environment specified in Section 1-2 and mount immediately.
- 3.Surface mount conditions.
- The following soldering condition are recommended to ensure device quality.

3-1.Soldering.

(1) Convection reflow or IR/Convection. (one-time soldering or two-time soldering)

• Temperature and period :

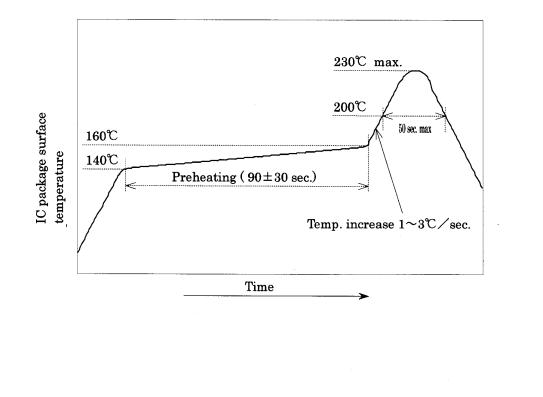
Peak temperature of 230°C max.

Above 200°C for 50 sec. max.

Preheat temperature of $140 \sim 160^{\circ}$ for 90 ± 30 sec.

Temperature increase rate of $1 \sim 3^{\circ} C / sec$.

- Measuring point : IC package surface.
- Temperature profile :

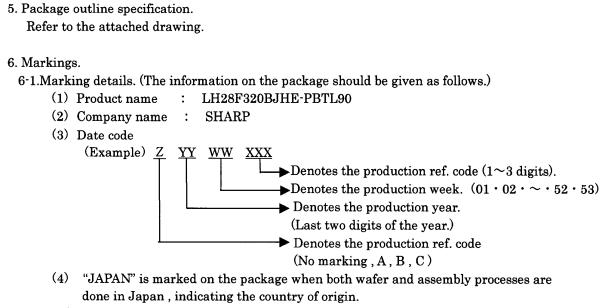


(2) Manual soldering (soldering iron) (one-time soldering only)

Soldering iron should only touch the IC's outer leads.

- Temperature and period :
 - $350\%\,$ max. for 3 sec. / pin max., or $260\%\,$ max. for 10 sec. / pin max.
 - (Soldering iron should only touch the IC's outer leads.) $% \label{eq:soldering}$
- $\boldsymbol{\cdot}$ Measuring point : Soldering iron tip.
- 4. Condition for removal of residual flax.
- (1) Ultrasonic washing power: 25 watts / liter max.
- (2) Washing time : Total 1 minute max.
- (3) Solvent temperature $:15 \sim 40^{\circ}$

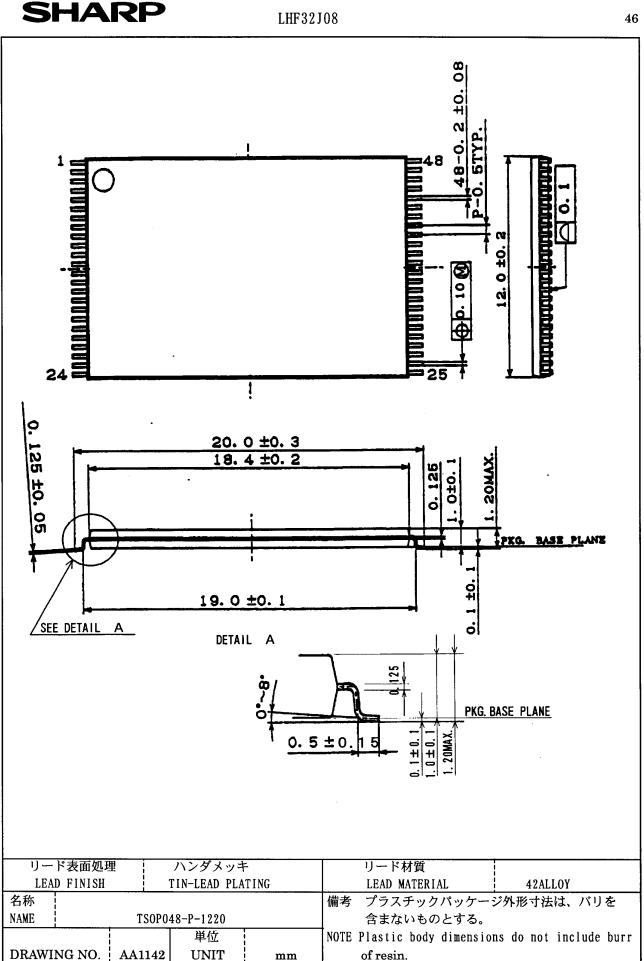


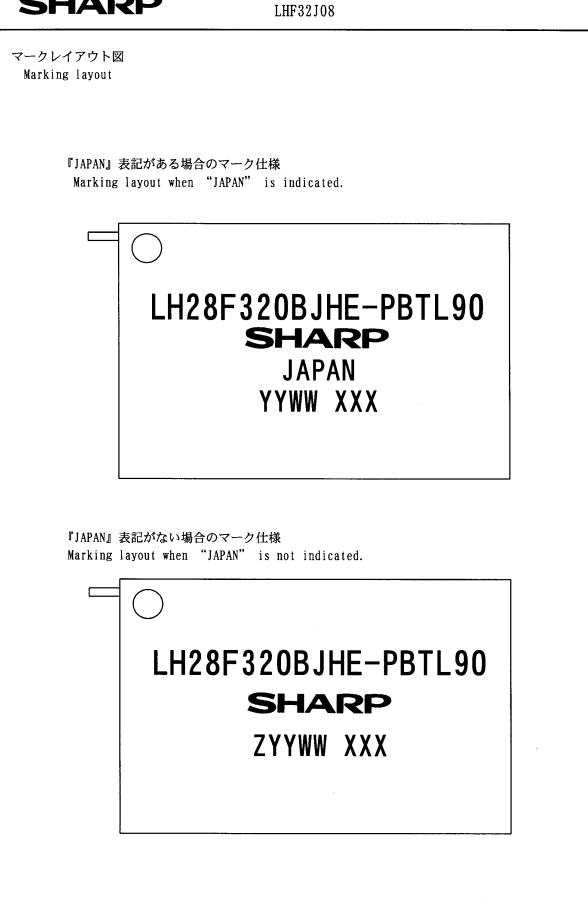


6-2.Marking layout.

The layout is shown in the attached drawing.

(However, this layout does not specify the size of the marking character and marking position.)





7.Packing Specifications (Dry packing for surface mount packages.) 7-1.Packing materials.

Material name	Material specifications	Purpose
Inner carton	Cardboard (500 devices / inner carton max.)	Packing the devices. (10 trays / inner carton)
Tray	Conductive plastic (50 devices / tray)	Securing the devices.
Upper cover tray	Conductive plastic (1 tray / inner carton)	Securing the devices.
Laminated aluminum bag	Aluminum polyethylene	Keeping the devices dry.
Desiccant	Silica gel	Keeping the devices dry.
Label	Paper	Indicates part number, quantity, and packed date.
PP band	Polypropylene (3 pcs. / inner carton)	Securing the devices.
Outer carton	Cardboard (2000 devices / outer carton max.)	Outer packing.

(Devices must be placed on the tray in the same direction.)

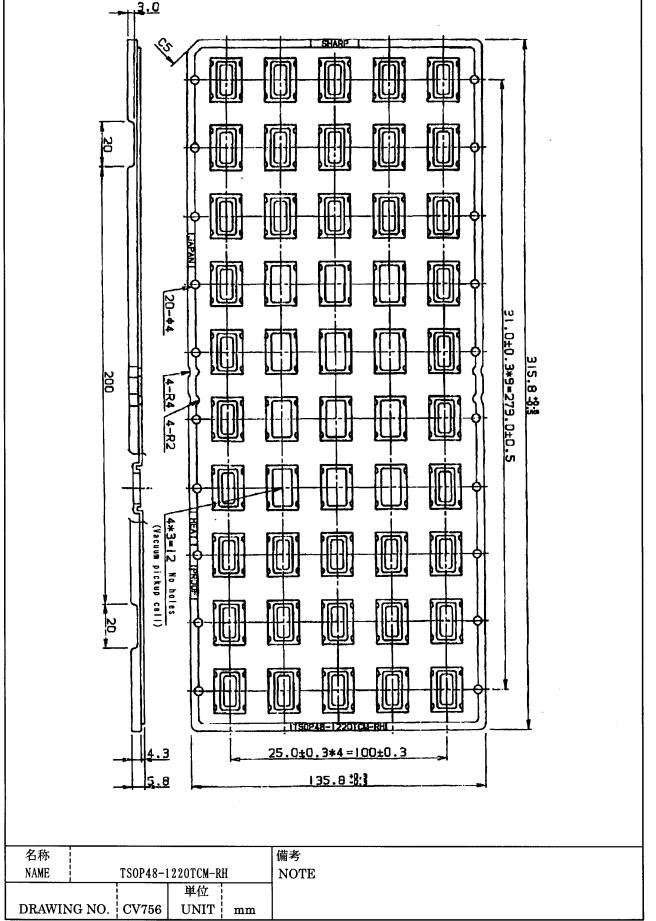
7-2.Outline dimension of tray.

Refer to the attached drawing.

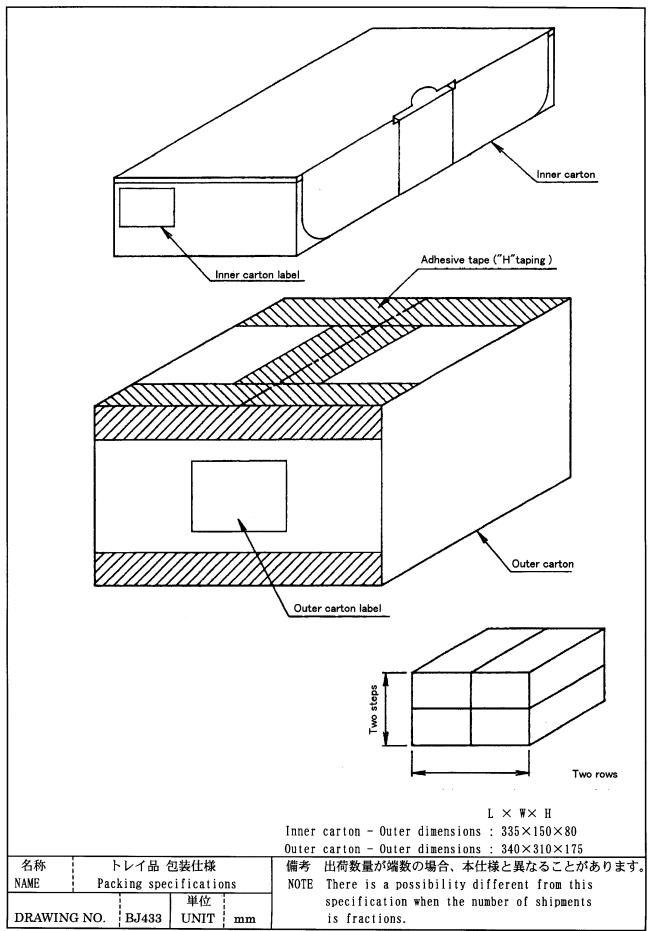
7-3.Outline dimension of carton.

Refer to the attached drawing.

- 8. Precautions for use.
 - (1) Opening must be done on an anti-ESD treated workbench. All workers must also have undergone anti-ESD treatment.
 - (2) The trays have undergone either conductive or anti-ESD treatment. If another tray is used, make sure it has also undergone conductive or anti-ESD treatment.
 - (3) The devices should be mounted the devices within one year of the date of delivery.







A-1 RECOMMENDED OPERATING CONDITIONS

A-1.1 At Device Power-Up

AC timing illustrated in Figure A-1 is recommended for the supply voltages and the control signals at device power-up. If the timing in the figure is ignored, the device may not operate correctly.

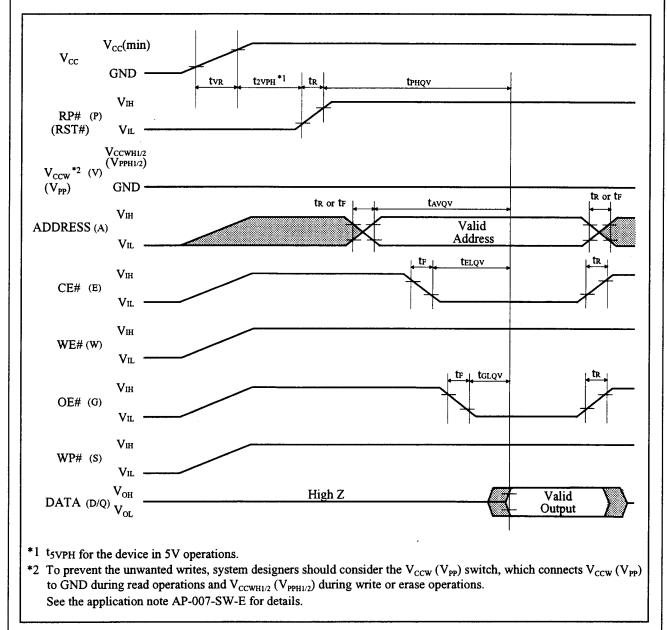


Figure A-1. AC Timing at Device Power-Up

For the AC specifications t_{VR} , t_R , t_F in the figure, refer to the next page. See the "ELECTRICAL SPECIFICATIONS" described in specifications for the supply voltage range, the operating temperature and the AC specifications not shown in the next page.

A-1.1.1 Rise and Fall Time

Symbol	Parameter	Notes	Min.	Max.	Unit
t _{VR}	V _{CC} Rise Time	1	0.5	30000	µs∕V
t _R	Input Signal Rise Time	1, 2		1	µs/V
t _F	Input Signal Fall Time	1, 2		1	μs/V

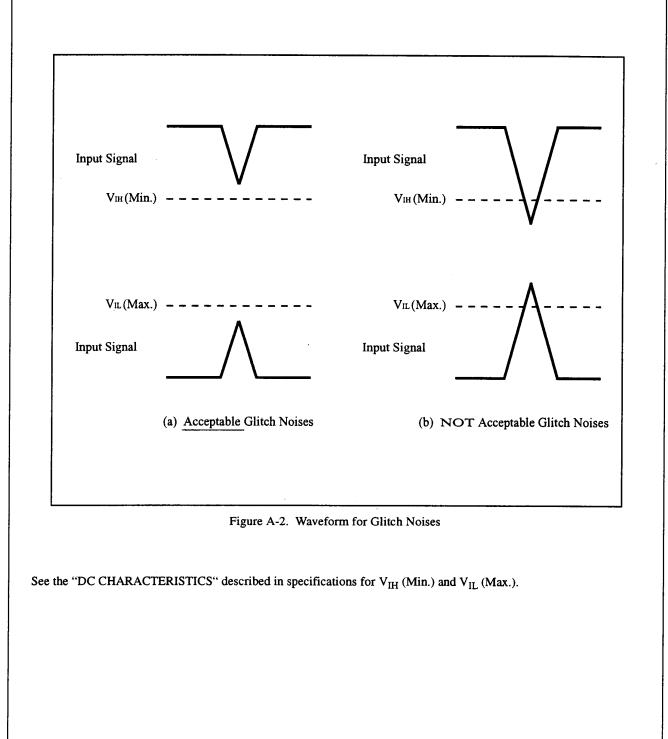
.

NOTES:

1. Sampled, not 100% tested. 2. This specification is applied for not only the device power-up but also the normal operations. $t_R(Max.)$ and $t_F(Max.)$ for RP# (RST#) are 50 μ s/V.

A-1.2 Glitch Noises

Do not input the glitch noises which are below V_{IH} (Min.) or above V_{IL} (Max.) on address, data, reset, and control signals, as shown in Figure A-2 (b). The acceptable glitch noises are illustrated in Figure A-2 (a).



A-2 RELATED DOCUMENT INFORMATION⁽¹⁾

Document No.	Document Name			
AP-001-SD-E	Flash Memory Family Software Drivers			
AP-006-PT-E	Data Protection Method of SHARP Flash Memory			
AP-007-SW-E	RP#, V _{PP} Electric Potential Switching Circuit			

NOTE:

1. International customers should contact their local SHARP or distribution sales office.

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